

# Geometric Constraint Lecture(Mar 21-23)

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## 1 Problem Categories

## 2 Five Questions

1. **Given graph  $G$ , characterize  $d$  for which  $(G, d)$  has a realization.**  
Here  $d$  are constraints, for example distance constraints.
2. **Given graph  $G$  and constraints  $d$ , provide a realization.**
3. **Given graph  $G$ , generically classify it into two categories:**
  - It has finite number of realizations.
    - One realization
    - Many realizations
  - It has infinite number of realizations.
4. **Given  $G$ , generically characterize the realization space.**
5. **Given nongeneric  $G$ , with fixed or restricted  $d$ , answer question 3 and 4. Give the classification and description of its realization space.**

## 3 Working on these Five Questions

### 3.1 Question 1

**Problem:**  $G$  is a complete distance graph, find  $\{d : (G, d) \text{ has a realization in } \mathbb{R}^k \text{ space}\}$ .

**Theorem:** Cayley-Menger conditions are the necessary and sufficient conditions that  $(G, d)$  has a realization in  $\mathbb{R}^k$  space.

### 3.2 Question 4

Question 1 and 4 are equivalent in the sense that if we understand one of them, we understand the other.

### 3.3 Question 3

**3.3.1 Leman's theorem:** A graph  $G$  generically has only finitely many solutions iff the following two conditions hold:

1.  $\forall \text{subgraph } S \subseteq G, 2|V_S| - |E_S| \geq 3$
2.  $2|V_G| - |E_G| = 3$

**3.3.1.1 General Leman theorem:** A graph  $G = (V, E)$  generically has at most finitely many solutions iff  $\exists \text{subgraph } G' = (V, E')$  with  $E' \subseteq E$  such that

1.  $\forall \text{subgraph } S \subseteq G', 2|V_S| - |E_S| \geq 3$
2.  $2|V_{G'}| - |E_{G'}| = 3$

#### 3.3.1.2 Definition of Generic

**Embedding:** can be understood in the following three ways:

- $(x_1y_1x_2y_2 \cdots x_ny_n) \subseteq \mathbb{R}^{2n}$
- $d_{\bar{G}} \subseteq \mathbb{R}^{|V_G|}$
- $\mathbb{R}^{2n} \setminus E_2$

**Given  $(G, d)$  o/w  $d_G$**

$$\{(x_1y_1x_2y_2 \cdots x_ny_n) : (x_v - x_w)^2 + (y_v - y_w)^2 = d_{vw}^2 \forall (v, w) \in E(G) \setminus E_2\}$$

$\leftrightarrow$

$$d_{\bar{G}} = \{(x_a - x_b)^2 + (y_a - y_b)^2, \cdots (a, b) \notin E(G), : (x_v - x_w)^2 + (y_v - y_w)^2 = d_{vw}^2 \forall (v, w) \in E(G)\}$$

**There is one to one map between these two sets.**

**Typical embedding:** An embedding  $Q$  of  $G$  is generic if

$\exists$  a small enough neighborhood of  $(d_{Q,G})$ ,  $d_{Q,G \pm \xi}$ , all realization of  $(G, d_{Q,G \pm \xi})$  are rigid  $\Leftrightarrow$   $Q$  is rigid.

**Alternately, One can define a small enough neighborhood of  $Q$  itself, Require all corresponding realization to be rigid for their corresponding distance values.**

### 3.3.1.3 Rigidity of Graph:

**Def1:** A Graph is rigid  $\exists$  a generic embedding that is rigid.

**Def2:** A Graph is rigid if all generic embeddings are rigid.

**These two definitions turn out to be equivalent and we will give the proof in the following notes.**

**Def3:** A Graph is globally rigid if it is rigid &

$\forall$  generic embedding  $Q$  with distances  $d_{G,Q}$ ,  $Q$  is the unique generic realization of  $(G, d_{G,Q})$

**Def4:** An embedding  $Q$  of  $G$  is rigid if

$\exists$  a small enough  $\mathbb{R}^{2n}$  - neighborhood  $Q_\xi$ , such that for  $\forall Q' \in Q_\xi$ ,  $Q'$  is a realization of  $(G, d_{Q,G}) \Leftrightarrow Q'$  is a rigid motion of  $Q$ .

**Def5:** An embedding  $Q$  of  $G$  is generic if

$\exists$  a small enough  $\mathbb{R}^{2n}$  - neighborhood  $Q_\xi$ , such that  $\forall Q' \in Q_\xi$ ,  $Q'$  is rigid  $\Leftrightarrow Q$  is rigid.

**Question:** Given a particular embedding  $Q \in \mathbb{R}^{2n}$  of  $G$ , Decide

1. is  $Q$  generic?
2. is  $Q$  rigid?

### 3.3.2 Jackson-Jordon theorem

**for  $d = 2$ ,  $G = (V, E)$  generically has an unique solution (globally rigid)  $\Leftrightarrow G$  is redundantly rigid & 3-connected or it is a triangle.**

**Redundantly rigid:** removal of any edge preserves rigidity of  $G$ .

### 3.3.2.1 Hendrickson's theorem

$G = (V, E)$  is globally rigid in  $d$  dimension  $\Leftrightarrow G$  is redundantly rigid for  $d$  dimension &  $(d + 1)$  connected.

$\Leftarrow$ proved

$\Rightarrow$ Conely disproved for  $d \geq 3$

$\Rightarrow$ proved (Jackson-Jordon theorem)

### 3.3.3 Owen's theorem

A graph is quadratically solvable  $\Leftrightarrow$  it is not 3-connected.

**Quadratically** solving: A consraint system  $(G, d)$  is quadratically solvable if it is triangularizable into quadratics.

$\Leftarrow$ proved

$\Rightarrow$ For planar graph, a graph is quadratically solvable  $\Rightarrow$  it is not 3-connected

$\Rightarrow$ For general graph, open problem