## Optimal Joining Of Compacted Cells

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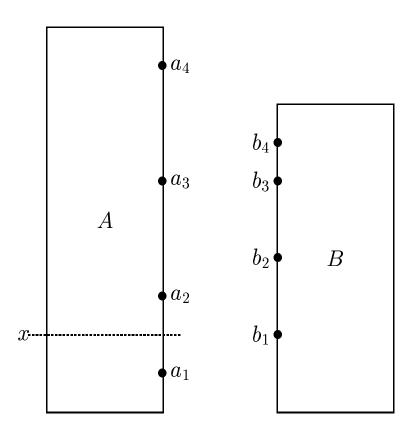
University of Florida

Siu-Wing Cheng

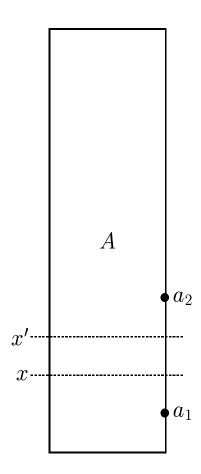
University of Minnesota

Sartaj Sahni

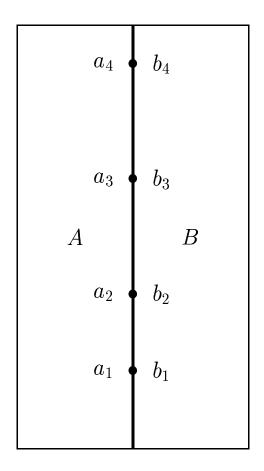
University of Florida



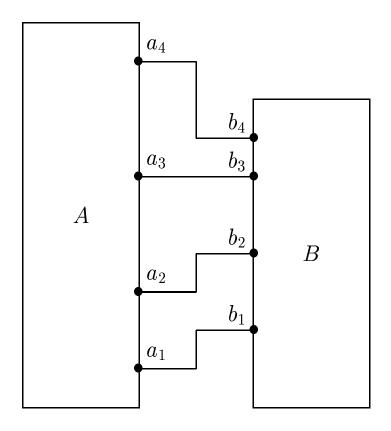
(a) Two adjacent cells



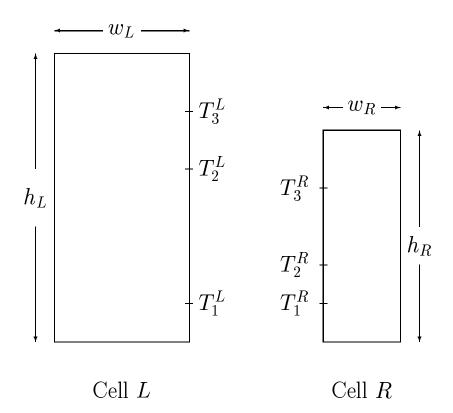
(b) stretching at x



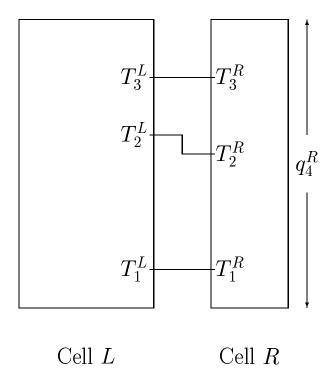
(c) joining by stretching



(d) joining by river routing



(a) Two compacted Cells



(b) After stretching and routing

```
procedure MinHeight(s);

/* Determine a minimum height legal stretching Q that can be routed using s tracks */

begin

q_0^L := 0; \ q_0^R := 0;

for i := 1 to s do begin

q_i^L := q_{i-1}^L + \Delta_{i-1}^L;

q_i^R := q_{i-1}^R + \Delta_{i-1}^R;

end;

for i := s + 1 to n do begin

q_i^L := q_{i-1}^L + \Delta_{i-1}^L;

q_i^R := q_{i-1}^L + \Delta_{i-1}^R;

if q_i^R < q_{i-s}^L + s then

q_i^R := q_{i-s}^L + s then

q_i^R := q_{i-s}^L + s then

q_i^L := q_{i-s}^R + s;

end;

end;
```

Figure 1: Procedure to find minimum height legal stretch

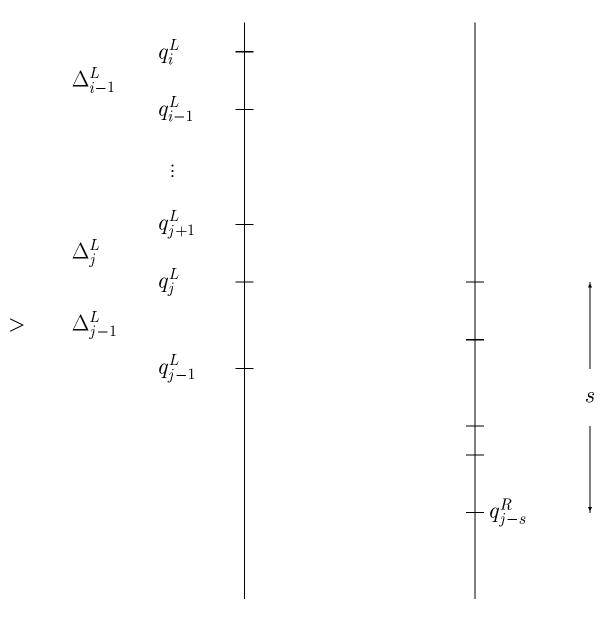


Figure 2: Relationship

```
procedure FixedLabeling(Q);

/* label some terminals as fixed. The labeled terminals satisfy conditions C1-C4 */
begin

Initialize a queue to be empty;

if q_n^L + \Delta_n^L \geq q_n^R + \Delta_n^R then label T_n^L fixed and add to the queue;

if q_n^R + \Delta_n^R \geq q_n^L + \Delta_n^L then label T_n^R fixed and add to the queue;

while queue not empty do begin

delete a terminal, say T_i^A, from the queue;

find least j such that q_{a+1}^A - q_a^A = \Delta_a^A, j \leq a < i;

label T_a^A, j \leq a < i, fixed;

if j > s then begin

let B = \{L, R\} - A;

add T_{j-s}^B to the queue;

label T_{j-s}^B, ..., T_j^B fixed;

end;

end;
end;
```

Figure 3: Procedure to label fixed terminals

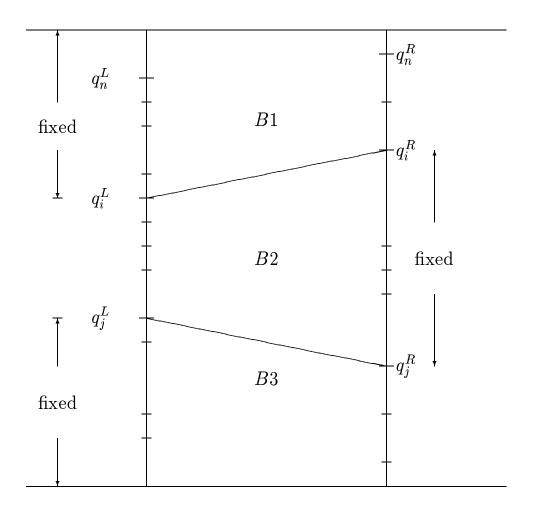


Figure 4: Block Decomposition

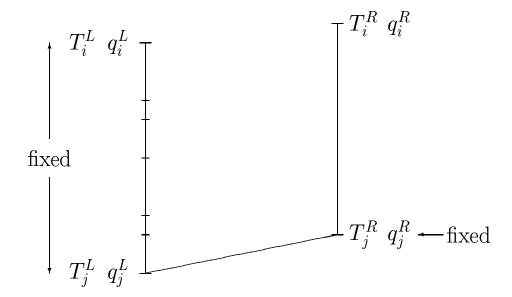


Figure 5: A block with left terminals fixed

```
procedure LeftMaxL(i,j);
      /* Determine the minimum maximum wire length for
            a block with fixed terminals on the left side. The
            terminals in the block are T_a^A, A \in \{L, R\}, j \le a \le i */
begin
      if i = n then q_{n+1}^R = MinHeight;
      high := q_{i+1}^R;
      MaxL := 0; /* length of longest wire */
      for a := i downto j + 1 do
            if q_a^L + MaxL < q_a^R
then begin /* wire a is longer than MaxL */
                   \begin{aligned} MaxL &:= q_a^R - q_a^L; \\ high &:= min\{high - \Delta_a^R, q_{a+s}^L - s\}; \end{aligned}
             end
            else begin /* wire a may be longer than MaxL */
                  \begin{array}{l} \mathbf{d}_{a} := \min \{q_{a}^{L} + MaxL, q_{a+1}^{R} - \Delta_{a}^{R}, q_{a+s}^{L} - s\}; \\ high := \min \{high - \Delta_{a}^{R}, q_{a+s}^{L} - s\}; \\ \mathbf{if} \ q_{a}^{L} - q_{a}^{R} > MaxL \\ \mathbf{then} \ q_{a}^{R} := \min \{high, q_{a+1}^{R} - \Delta_{a}^{R}, q_{a}^{R} + \lceil (q_{a}^{L} - q_{a}^{R} - MaxL)/2 \rceil \}; \\ MaxL := \max \{MaxL, q_{a}^{L} - q_{a}^{R}\}; \end{array}
             end;
end;
```

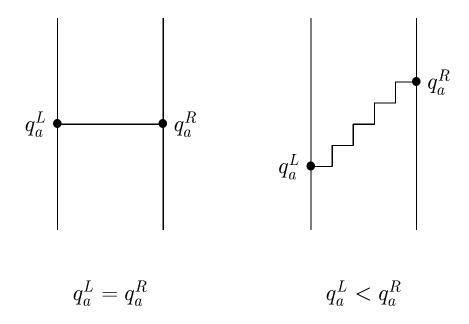
Figure 6: Determine minimum maximum length

```
procedure LeftRelocate(i,j);

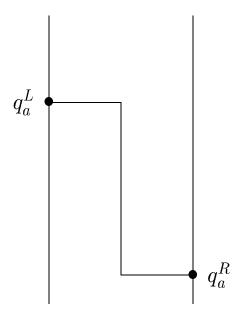
/* Relocate the right terminals of a block whose left terminals are fixed. The maximum wire length is to be MaxL and the terminals are T_a^A, A \in \{L, R\}, j < a \le i\} */
begin

for a := i downto j + 1 do
q_a^R := min\{q_a^L + MaxL, q_{a+1}^R - \Delta_a^R, q_{a+s}^L - s\}
end;
```

Figure 7: Relocate the right moveable terminals



(a) Positive terminals



$$q_a^L > q_a^R$$

(b) Negative terminals

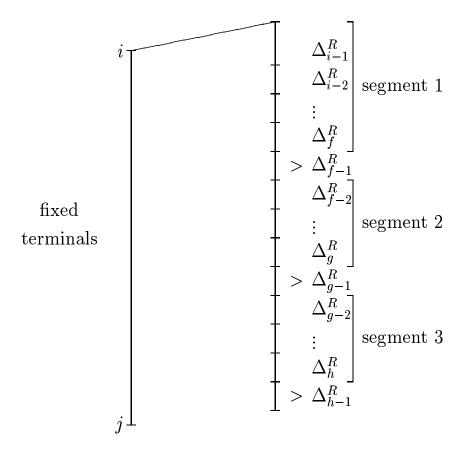


Figure 8: Segments

```
procedure Update(i,j);
   /* Use Offsets to compute q_a^{Rl}s */
begin
Initialize an empty stack;
AddValue := 0; EndIndex := i
for a := j to i do
   begin
       p := Offset(a);
       while p \ll nil do /* get offsets from the list */
          get(x, SegBegin) from node p;
          p := \text{next node on the offset list};
          AddToStack((AddValue, EndIndex));
          AddValue := AddValue + x;
          EndIndex := SegBegin;
       q_a^R := q_a^R + AddValue;
      if (a = EndIndex) and (a < i)
       then DeleteFromStack((AddValue, EndIndex));
   end;
end;
```

Figure 9: Second Pass of total wire length algorithm