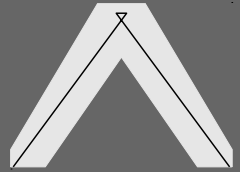
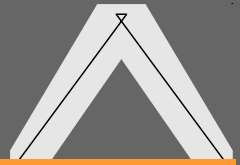


Goals and Outline



- (1) SLEFEs: Enclosing Functions (and SLEVEs for curves and surfaces)
- (2) *Mid-structures: Quantitatively* Coupling Curved and pw Linear Geometry
- (3) *Constrained Design*: One-sided fitting

Midstructures

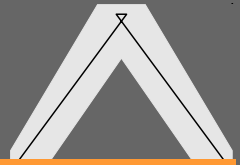


$$\begin{aligned}\underline{\bar{x}}(t) &:= (\bar{x}(t) + \underline{x}(t))/2 = \frac{1}{2} \left(\sum_{\mu=0}^m \tilde{x}_{\mu} \mathbf{h}_{\mu}^m(t) + \sum_{\mu=0}^m \tilde{x}_{\mu} \mathbf{h}_{\mu}^m(t) \right) \\ &= x_0(1-t) + x_d t + \sum_{\mu=1}^m \sum_{\nu=1}^{d-1} \mathcal{F}_{\nu} x \frac{a[d, m, +, \nu, \mu] + a[d, m, -, \nu, \mu]}{2} \mathbf{h}_{\mu}^m(t) \\ &= Lx(t) + \mathcal{F}x \cdot \underline{\bar{a}} \cdot \mathbf{h}(t), \quad \underline{\bar{a}} \in \mathbb{R}^{d+1 \times m+1}.\end{aligned}$$

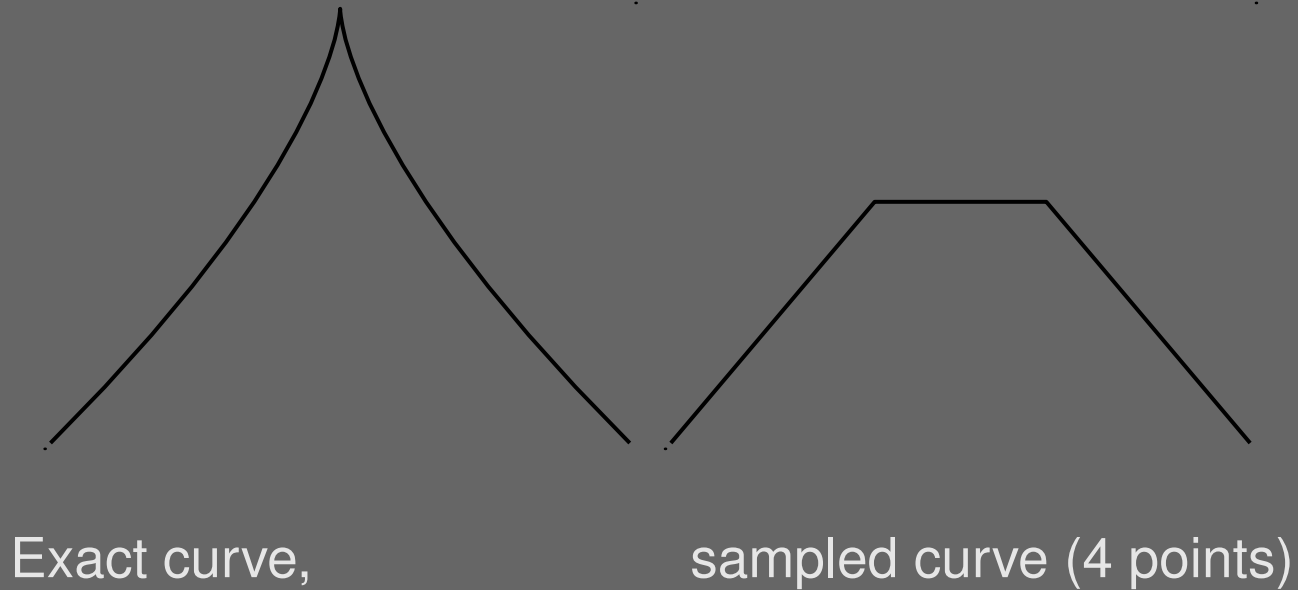
- Modification for Bézier: endpoints are x_0 and x_d .
- Similar to center curve for interval splines.
- Well-defined for vector-valued curve or surface.

Why is it useful?

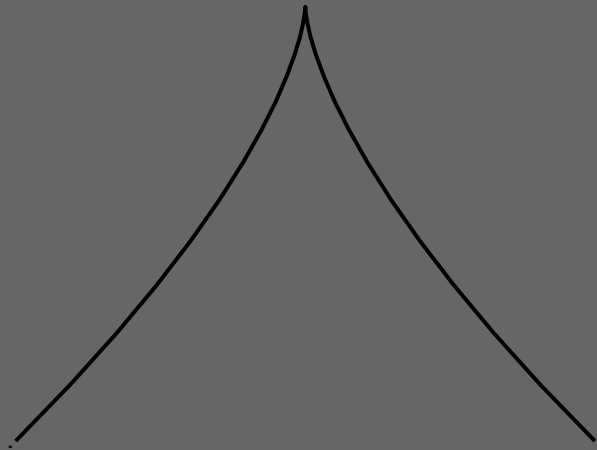
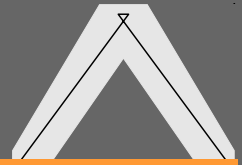
Application: Evaluation



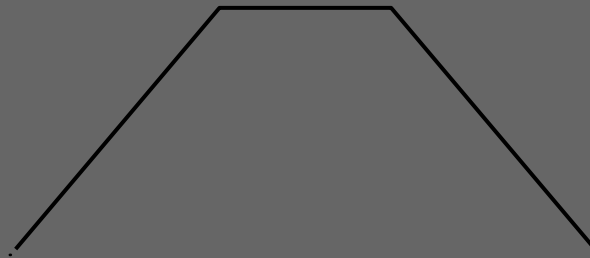
compare to sampling:



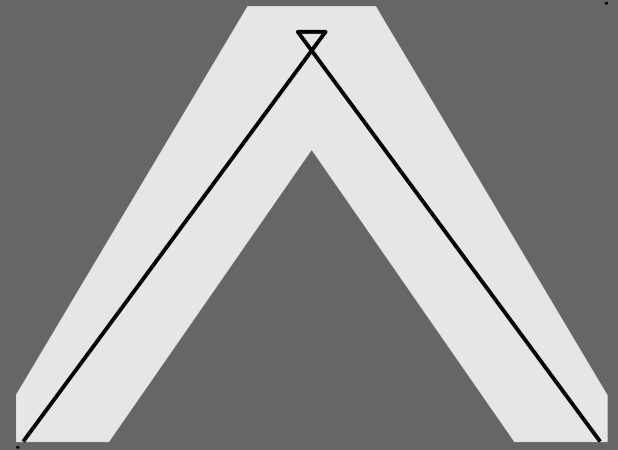
Application: Mid-Path



Exact curve,



sampled curve,

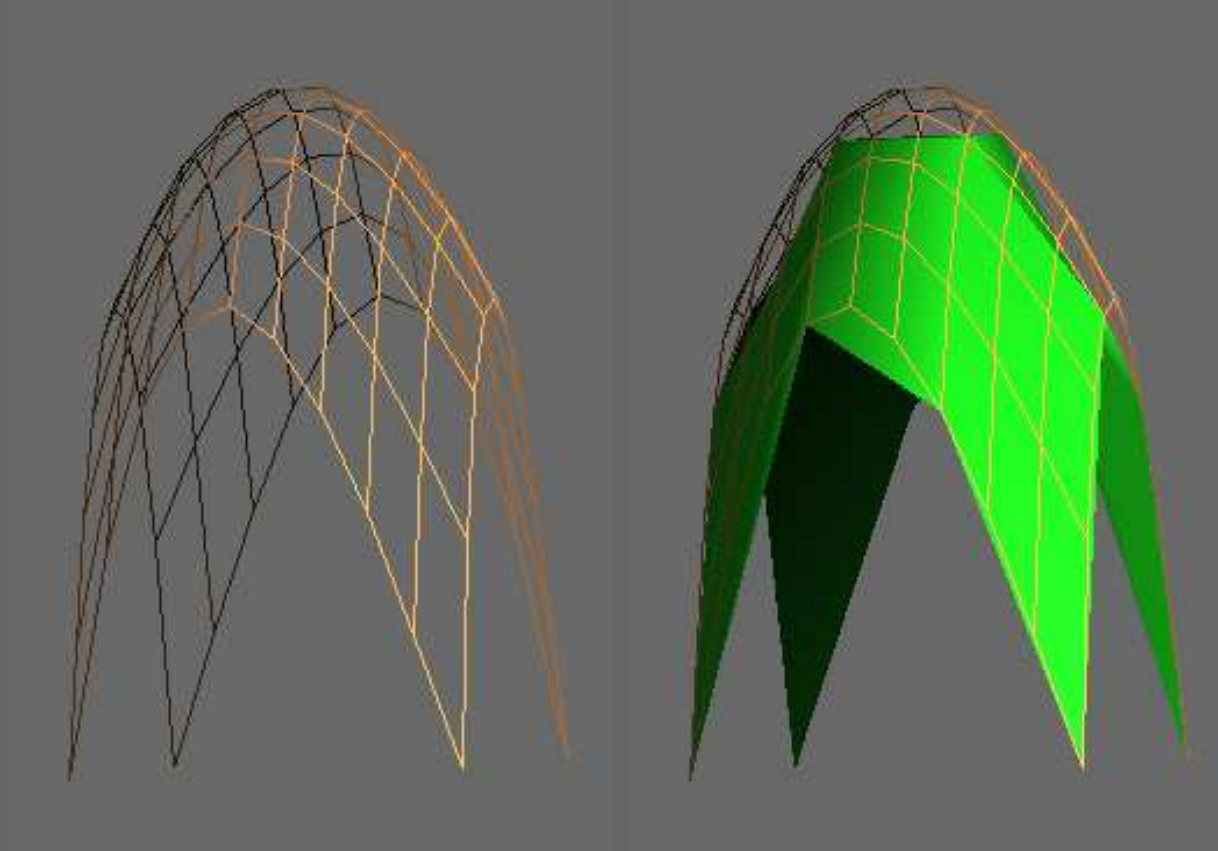
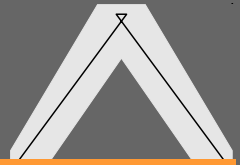


mid-path

max-norm approximation,
lower cost* than evaluation,
error bounds available.

$$\text{Mid-path} = (\underline{x} + \bar{x}) / 2.$$

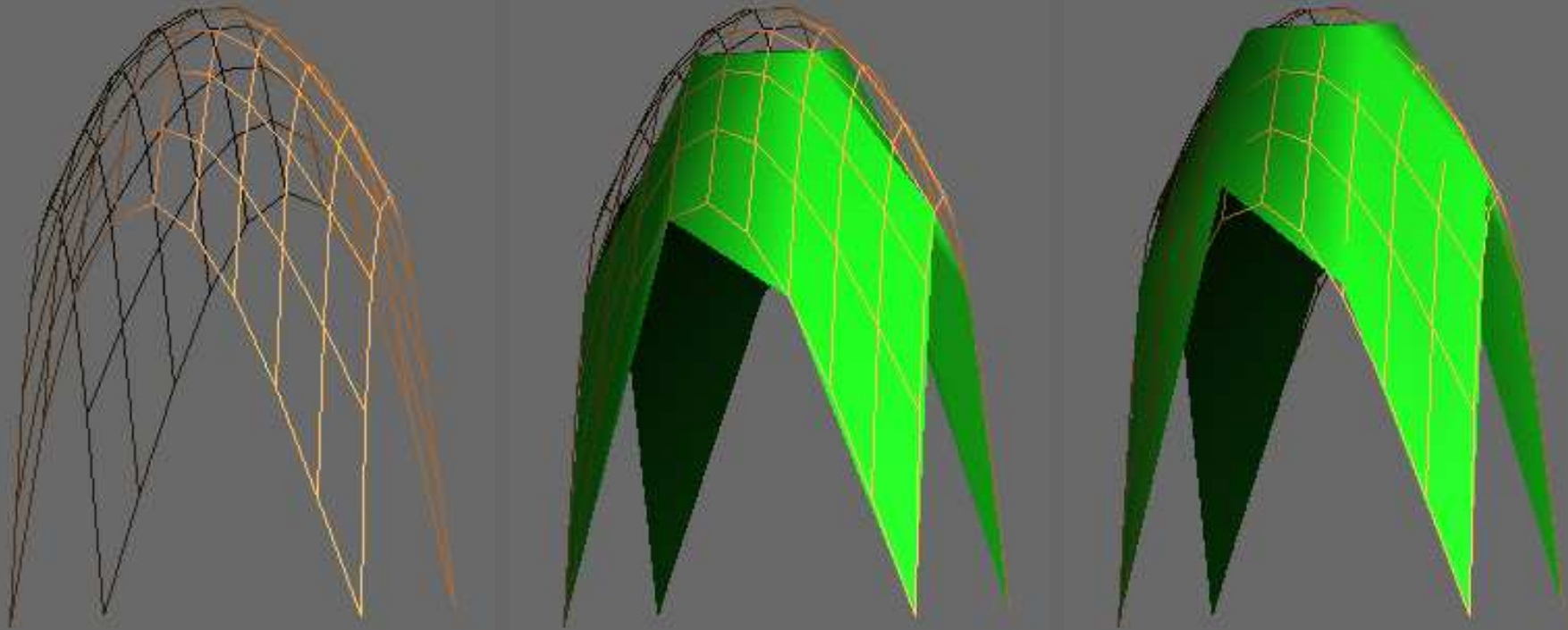
Mid-Patch: Approximate Rendering



finely sampled,

coarsely sampled (green): budget=16 points

Mid-Patch: Approximate Rendering



max-norm approximation,
lower cost than evaluation,
free error bounds.

(Bézier: boundary is defined by control points of the boundary only.)

Max-norm approximation



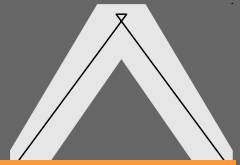
Mid-structures equi-oscillate.

Why not Chebyshev economization?

Why not Remez algorithm?

One-sided approximation!

Midpath Control Structures (Inversion)



If *number of segments = number of control points*, e.g. $m = d$
then $\underline{\bar{a}}$ is invertible*!

Reverse the midpath *coefficient* computation:

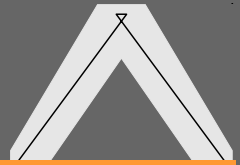
$$\underline{\bar{x}} - Lx = \mathcal{F}x \quad \underline{\bar{a}}.$$

solve for $\mathcal{F}x$ and reconstruct **the** function

$$x(t) = Lx(t) + \mathbf{a}(t) \cdot \mathcal{F}x$$

mid-struct and control-polytope = spline coefficients in *different bases*!

Midpath Control Structures (Inversion)



Any broken line \bar{x} (plus boundary conditions Lx)
has an associated x
(Thm:) that lies in the same flat*.

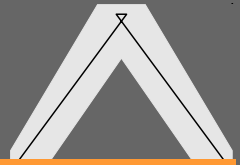
Examples



spline curve

box-spline

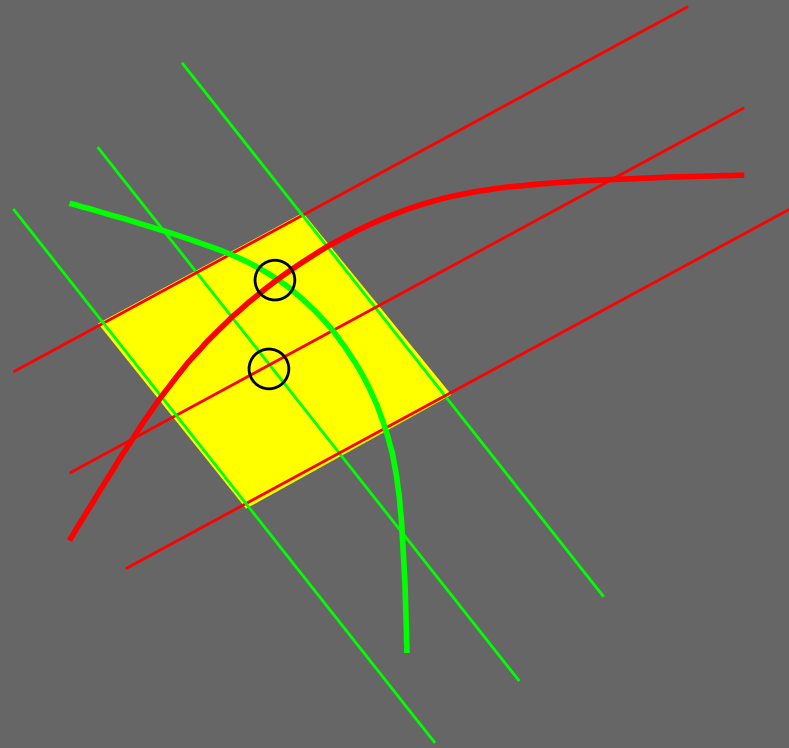
Application: Intersection



Robustness: *pw-linear*

Consistency: *intersect mid-structures*

2D: error at most $2 \min\{\epsilon, \tilde{\epsilon}\}$ (false positive, false negative)



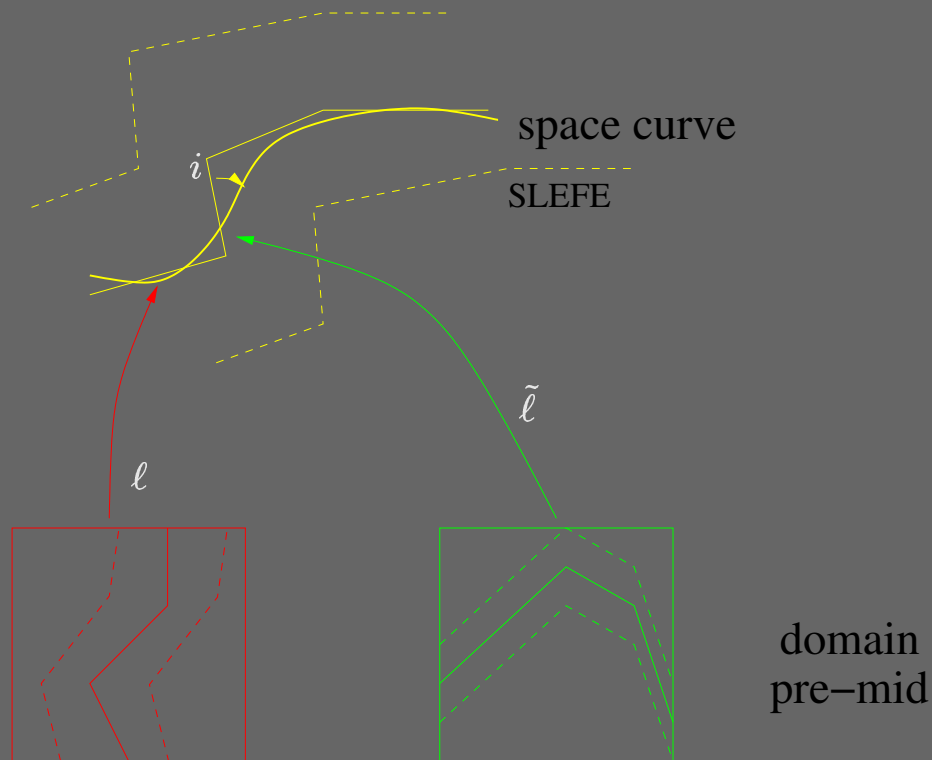
(P & Wu 2003): SLEFE construction to match a given ϵ .

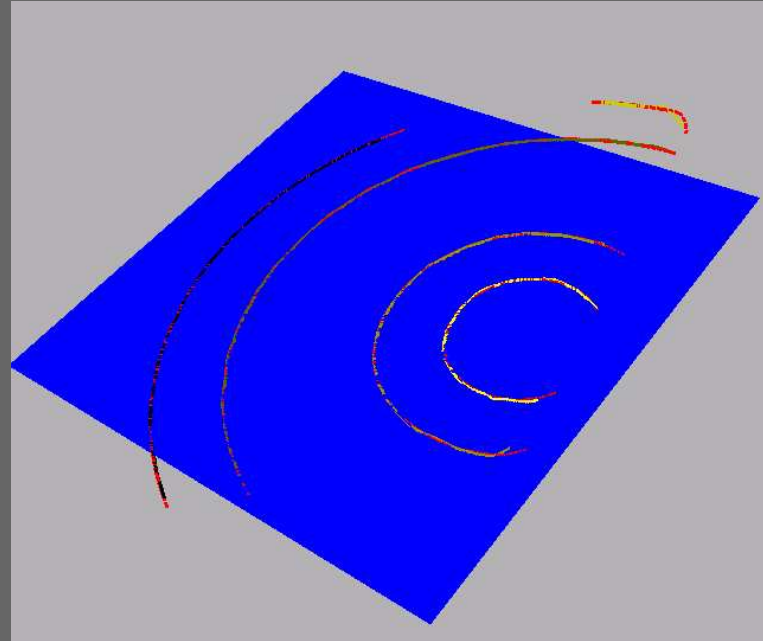
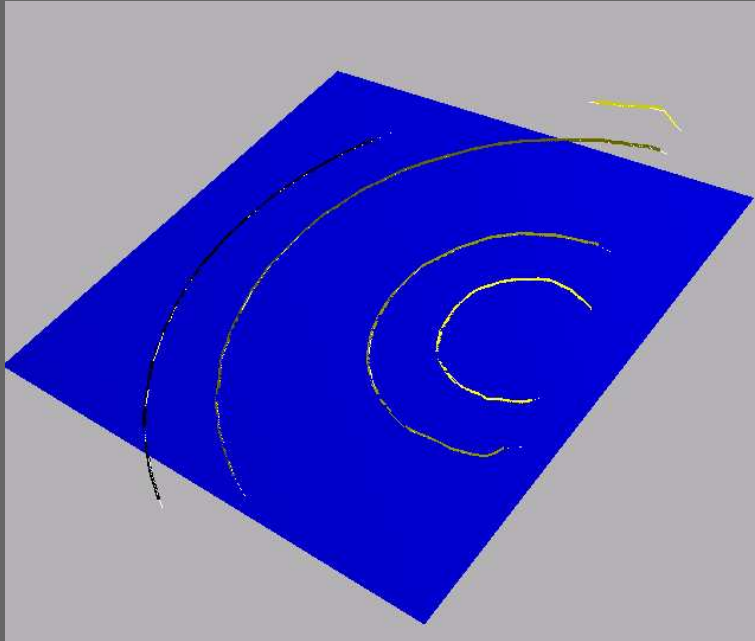
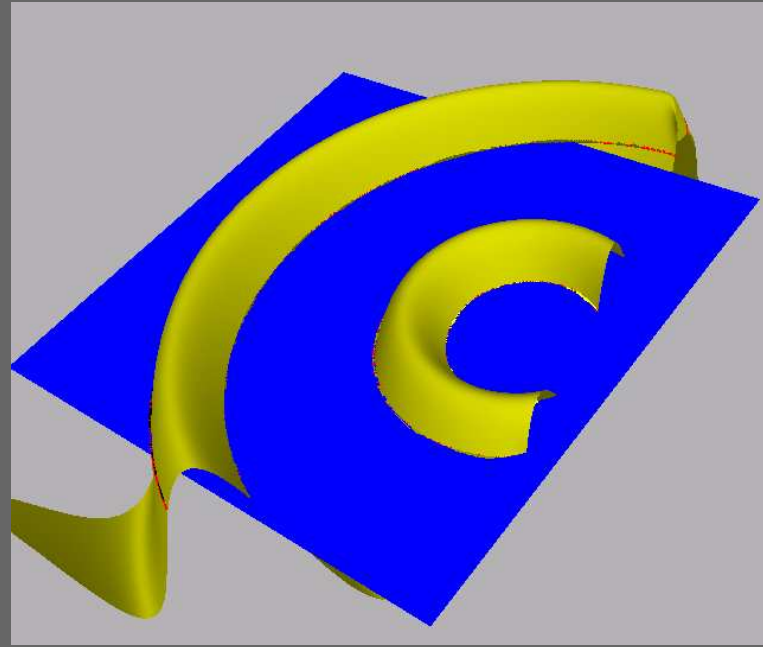
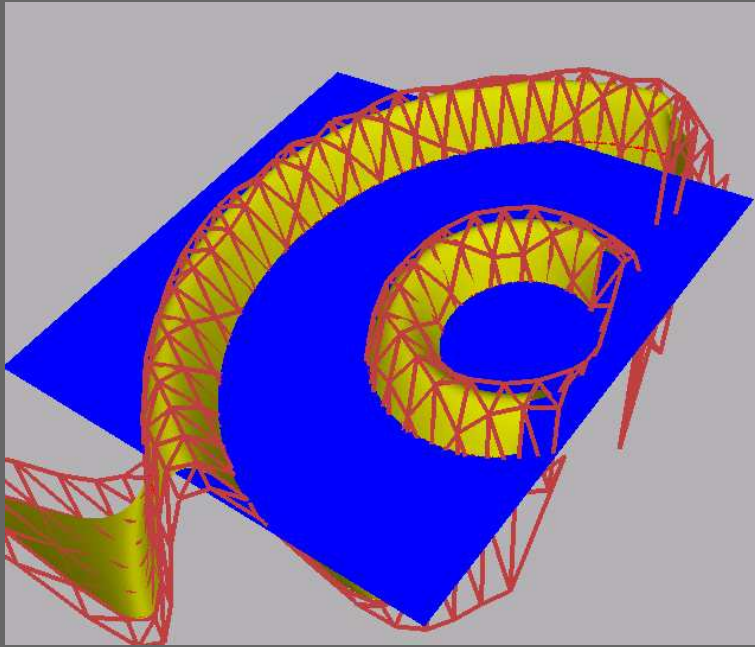
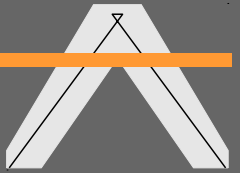
Application Intersection

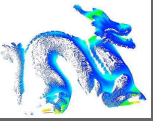


Robustness
Consistency

intersect pw-linear mid-structures 3D, 2 variables:
interpret intersection of $\underline{\bar{x}}$ and $\tilde{\underline{x}}$ as mid of a space curve.
Unique representer: $i(\ell(u)) = i(m) = i(\tilde{\ell}(\tilde{u}))$



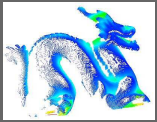




Goals and Outline

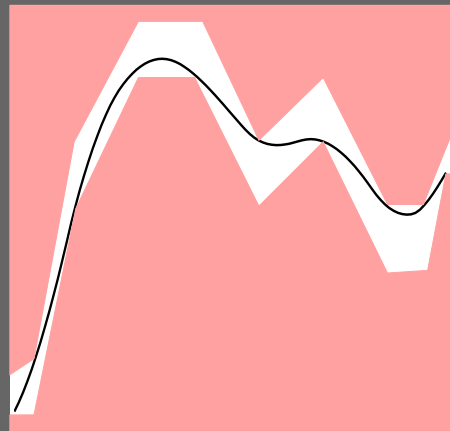
- (1) SLEFEs: Enclosing Functions
- (2) Mid-structures: Quantitatively Coupling Curved and pw Linear Geometry
- (3) Constrained Design: *One-sided fitting*





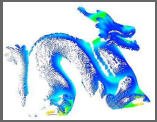
Inverse problems: Channel

Given two locally non-intersecting input polygons, construct a spline that stays between and consists of a small number of pieces.



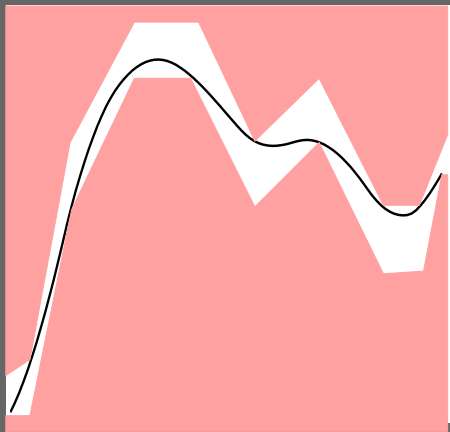
Channel

Continuous optimization problem!

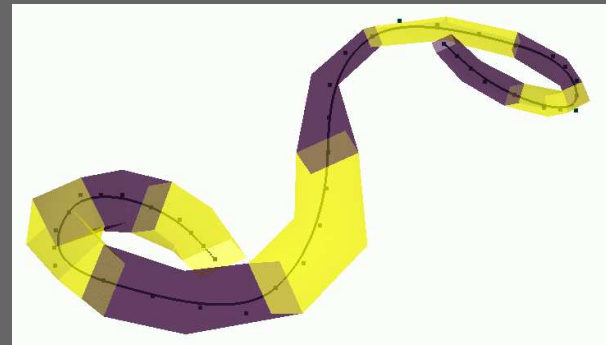


Inverse problems: Channel

Given two locally non-intersecting input polygons, construct a spline that stays between and consists of a small number of pieces.



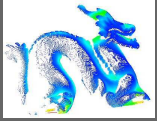
Channel



Channel 3D

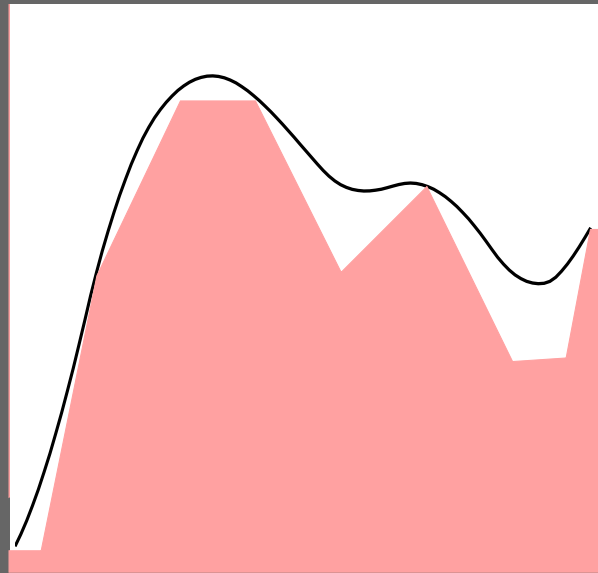
Solve for spline coefficients such that SLEFE does not intersect.

Idea: fit a SLEFE into the channel = *linearize and discretize* – a linear program!

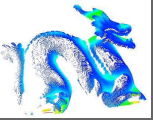


Inverse problems: Cover

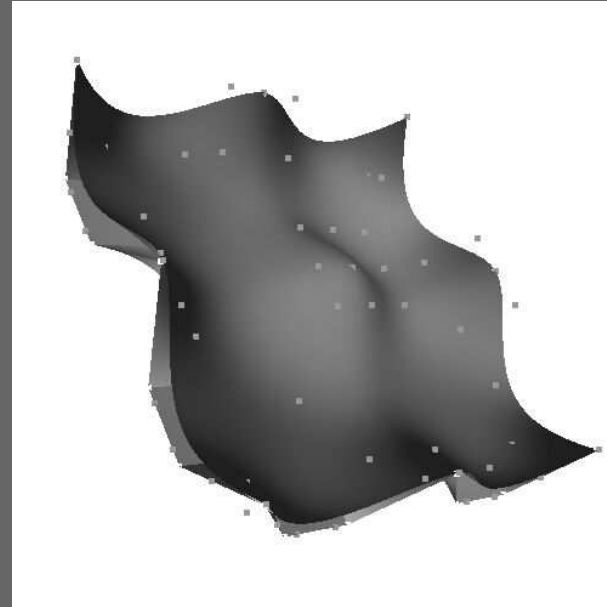
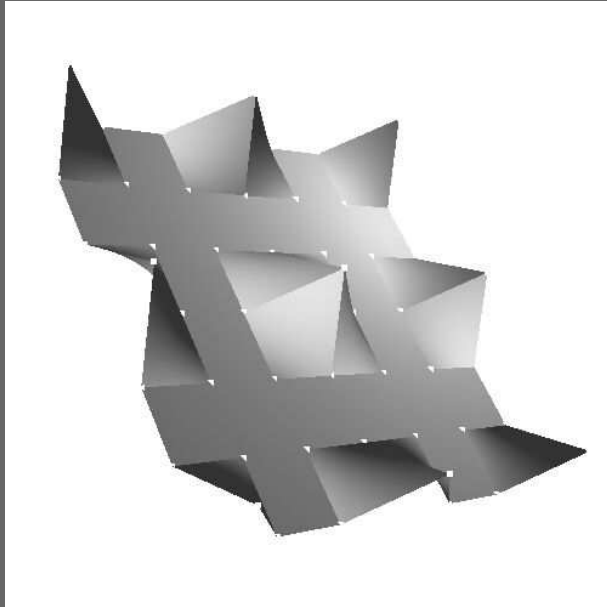
Given a barrier, construct x of a given degree and smoothness to stay close to and entirely above.

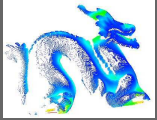


Motivation: bridge to computational geometry algorithms;
determinate assembly, simplification for intersection testing.



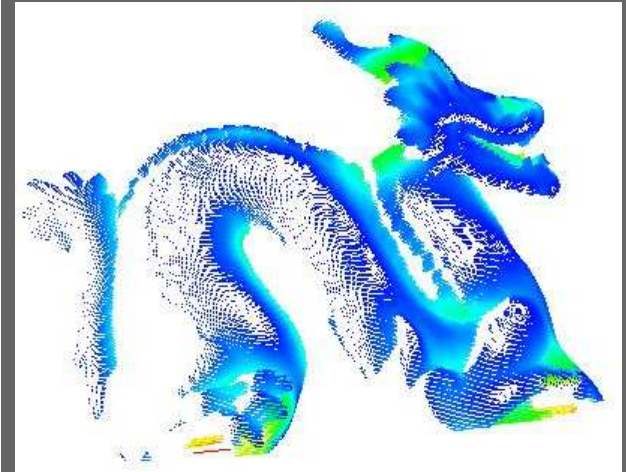
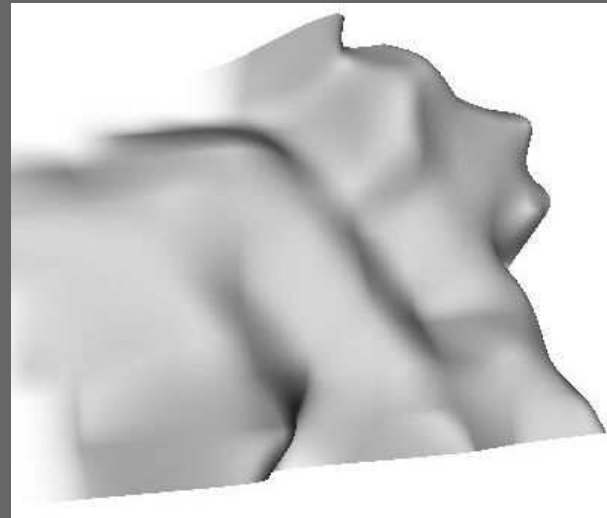
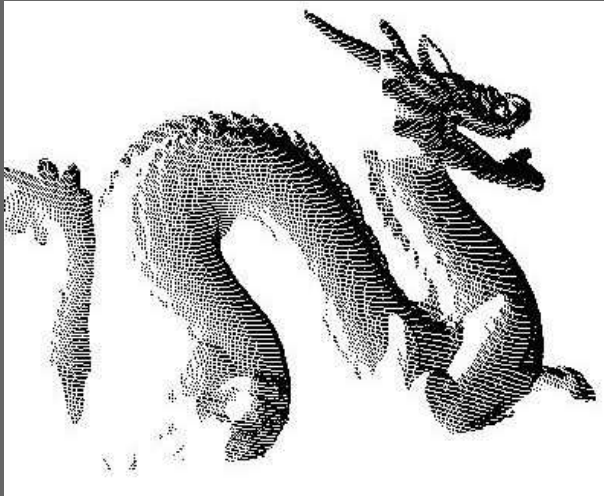
Bilinear cover



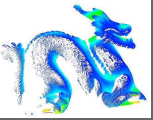


Range data fitting

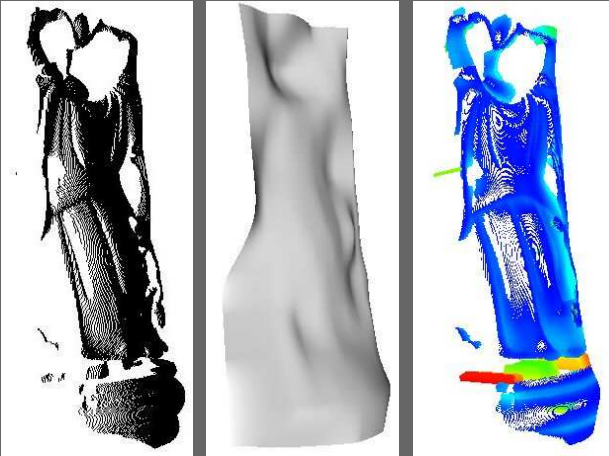
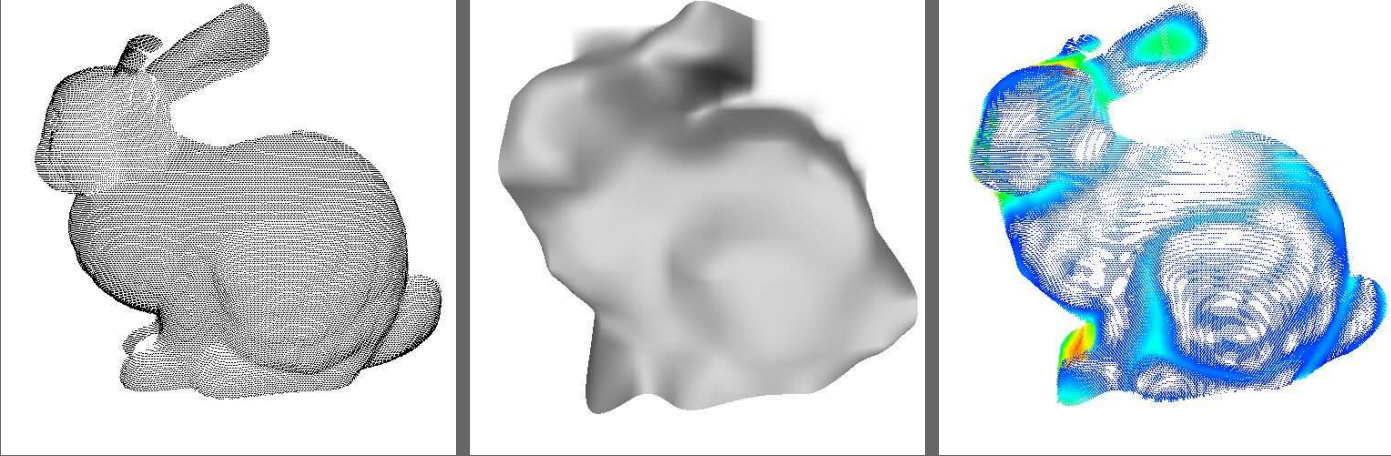
Left: Input range data (Dragon front view, 39526 points)
Middle: 10×10 bicubic spline
Right: offset distance from the range data.

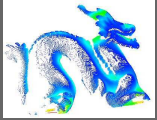


(orig) 4 min, (midpath fit + offset) 45 sec

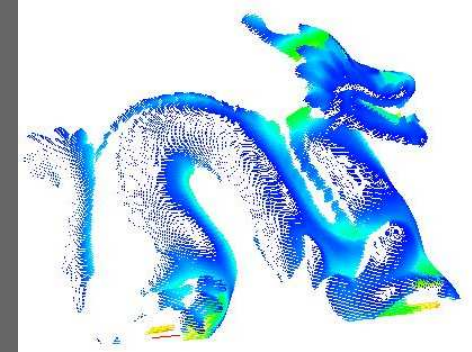
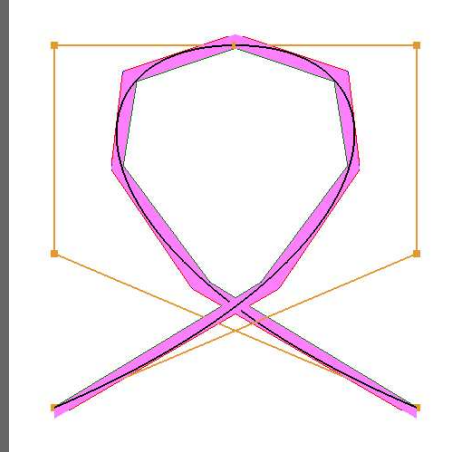
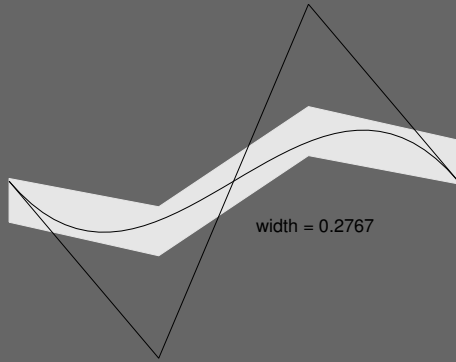


Range data fitting





Questions?



(Nairn, P, Lutterkort 1998; Reif) Sharp, quantitative bounds on the distance ...
(Lutterkort & P 2000): Optimized Refinable Enclosures of Multivariate Polynomial Pieces
(Myles & P 2003): Threading splines through 3D channels
(P & Wu 2003): SLEVEs for planar spline curves

