

Geometric Constraint Lecture

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Jan 31, 2006

1 Demo 1 from Solidwork software: Hexagon with three diagonals in 2D

Figure 1 is a Hexagon with three diagonals in 2D and the whole graph is rigid(wellconstrained, dense). The minimal rigid subgraph (wellconstrained subgraph, dense subgraph) is the whole graph. (Note: trival rigid subgraph, such as two vertices and one distance between them in 2D, is not considered!) So, the only DR-plan for this example is shown in the Figure 2.

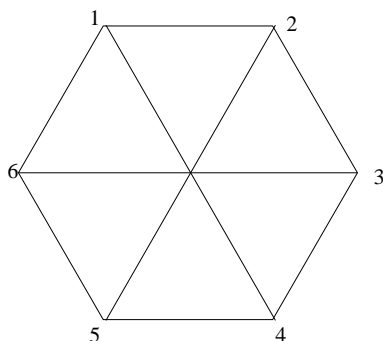


Figure 1: Hexagon with three diagonals in 2D

Figure 3 has a DR-plan shown in Figure 4. The fan-in (the maximum number of the children of a cluster in the DR-plan) is 3 in Figure 4 while the fan-in in Figure 2 is 6, so intuitively hexagon is more difficult to solve. Please recall the difference between Computational Geometry and Geometric Constraint Solving that we have introduced before: in Computational Geometry, the size of the minimum equation/inequality system we need to solve is bound by some constant, while in Geometric Constraint Solving, the

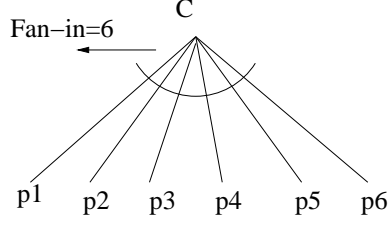


Figure 2: The only DR-plan for Hexagon with three diagonals in 2D

size can not be bound by some constant, for example it can be only bound by the input size.

Question: Can we give a 2D distance constraint graph $G=(V,E)$ (the only type of constraint is distance constraint) in which the minimum of fan-in of all possible DR-plans is $\theta(|V|)$?

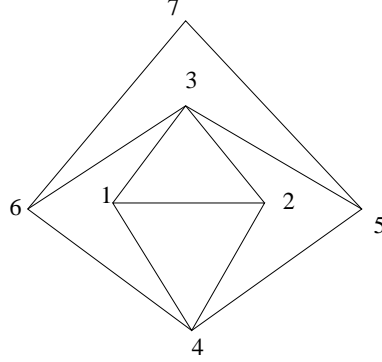


Figure 3: A Ruler Compass Constructable Example in 2D

2 Assembly example in Solidworks Software

Figure 5 is an assembly example in which $C1$ and $C2$ are two rigid bodies. $C1$ and $C2$ are assembled by the constraints 1378 and 2456 are coplanar and 13 and 24 are colinear.

For rigid 3D cluster $C1$ and $C2$: $DOF(C1) = DOF(C2) = 6$. The coplanarity will remove 3 DOF, the colinearity will remove 2 more extra DOF and the zero distance between 1 and 2 will remove 1 more extra DOF. So, the whole system is still rigid because the $DOF=6+6-3-2-1=6$.

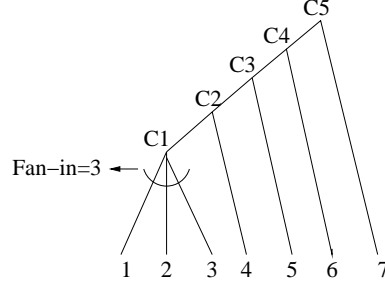


Figure 4: DR-plan for a RCC example

3 Demo of Frontier: solution navigation

In Figure 6, the point 1, 2 and 3 are the centers of three circles which are tangent to each other. Figure 7 is one of its DR-plan.

4 Banana in 3D

In Figure 8, there are two clusters C1 and C2 which are nontrivially rigid. By counting the DOF of the whole constraint graph, $2 \cdot 6 - 6 = 6$, so the whole graph should be rigid if the way counting DOF is correct! However, the whole system is either non-rigid and C1 and C2 can relatively rotate around hinge 13 if C1 and C2 provides a consistent distance for 13, or there is no real embedding if C1 and C2 provides a consistent distance for 13. This example also we can not check the rigidity of even a generic system by counting the DOF of its constraint graph.

In Figure 9, an extra distance between 24 can prevent the relative rotation between C1 and C2.

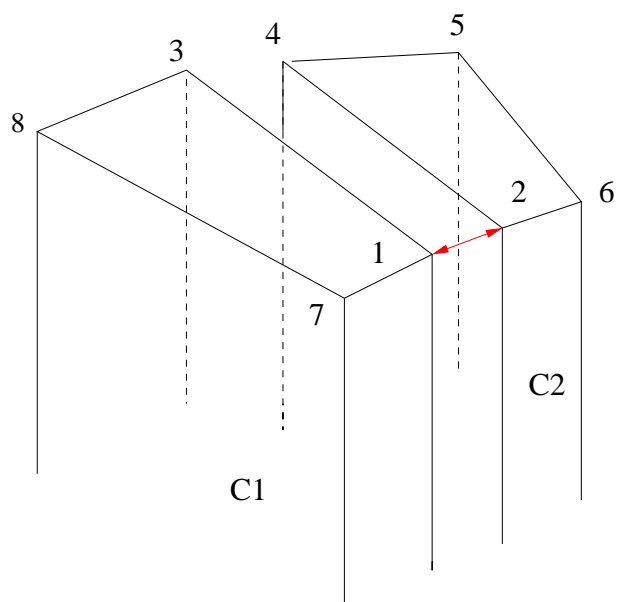


Figure 5: Assembly example of coplanarity and colinearity

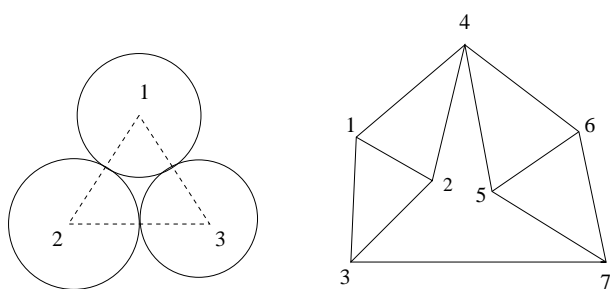


Figure 6: Frontier Demo

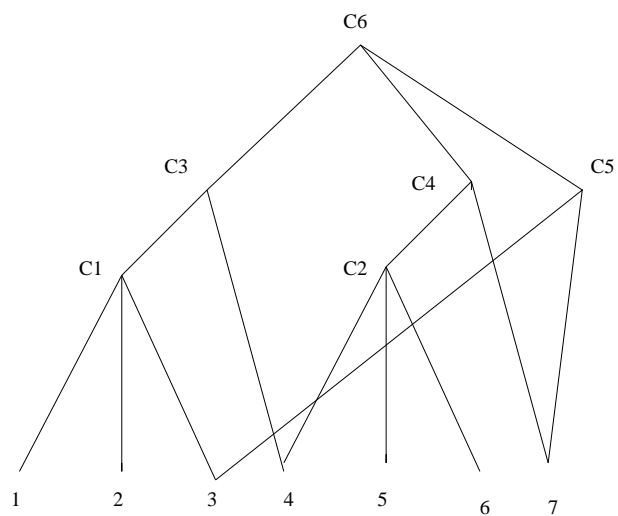


Figure 7: DR-Plan for the Frontier Demo

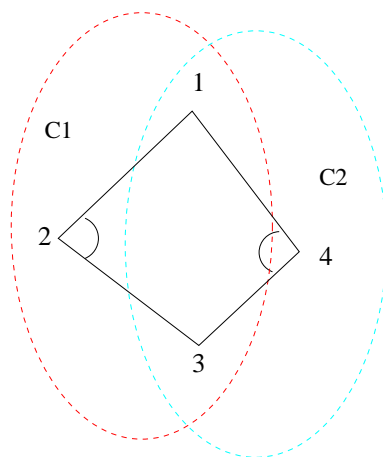


Figure 8: Banana example

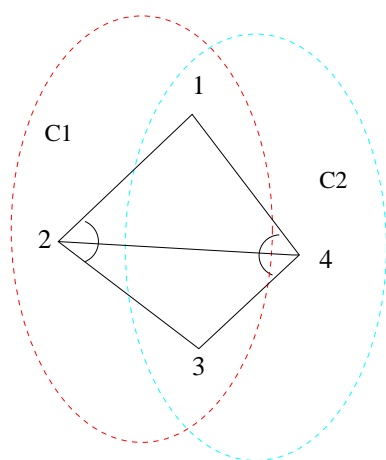


Figure 9: Banana example