

## Backtracking And Branch And Bound



## Subset & Permutation Problems

- Subset problem of size  $n$ .
  - Nonsystematic search of the space for the answer takes  $O(p2^n)$  time, where  $p$  is the time needed to evaluate each member of the solution space.
- Permutation problem of size  $n$ .
  - Nonsystematic search of the space for the answer takes  $O(pn!)$  time, where  $p$  is the time needed to evaluate each member of the solution space.
- Backtracking and branch and bound perform a systematic search; often taking much less time than taken by a nonsystematic search.

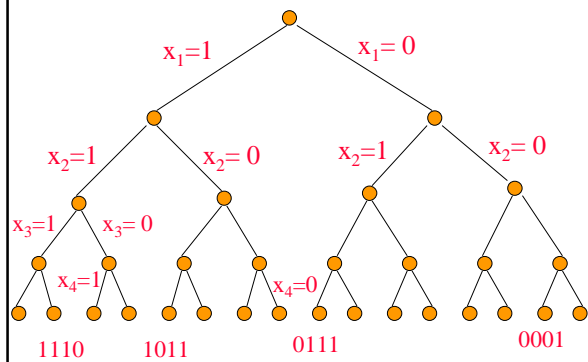
## Tree Organization Of Solution Space

- Set up a tree structure such that the leaves represent members of the solution space.
- For a size  $n$  subset problem, this tree structure has  $2^n$  leaves.
- For a size  $n$  permutation problem, this tree structure has  $n!$  leaves.
- The tree structure is too big to store in memory; it also takes too much time to create the tree structure.
- Portions of the tree structure are created by the backtracking and branch and bound algorithms as needed.

## Subset Problem

- Use a full binary tree that has  $2^n$  leaves.
- At level  $i$  the members of the solution space are partitioned by their  $x_i$  values.
- Members with  $x_i = 1$  are in the left subtree.
- Members with  $x_i = 0$  are in the right subtree.
- Could exchange roles of left and right subtree.

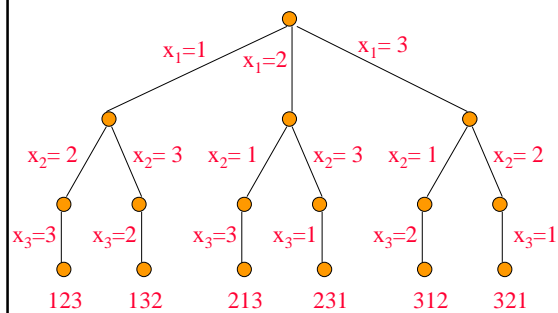
### Subset Tree For $n = 4$



### Permutation Problem

- Use a tree that has  $n!$  leaves.
- At level  $i$  the members of the solution space are partitioned by their  $x_i$  values.
- Members (if any) with  $x_i = 1$  are in the first subtree.
- Members (if any) with  $x_i = 2$  are in the next subtree.
- And so on.

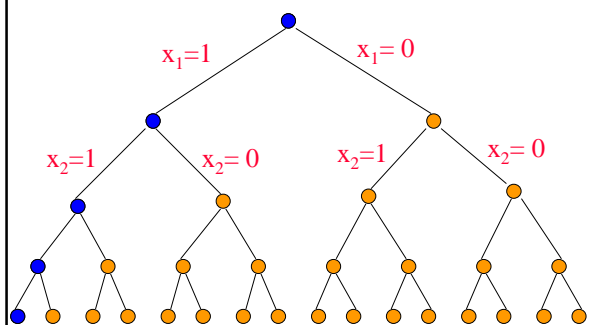
### Permutation Tree For $n = 3$



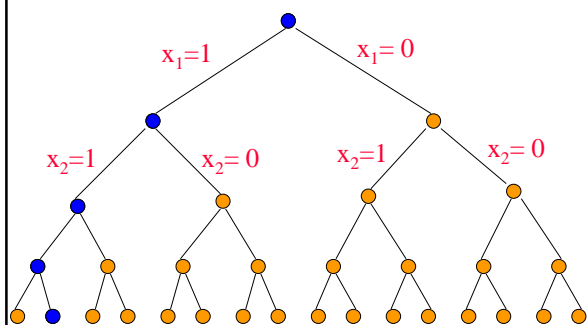
### Backtracking

- Search the solution space tree in a depth-first manner.
- May be done recursively or use a stack to retain the path from the root to the current node in the tree.
- The solution space tree exists only in your mind, not in the computer.

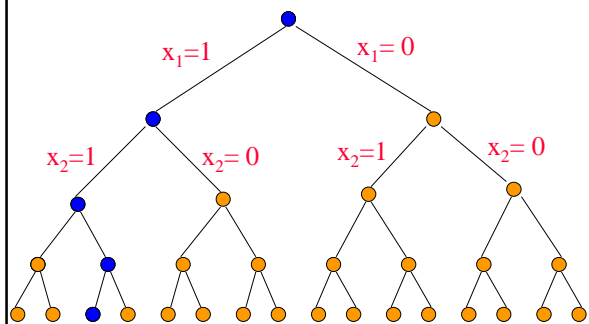
Backtracking Depth-First Search



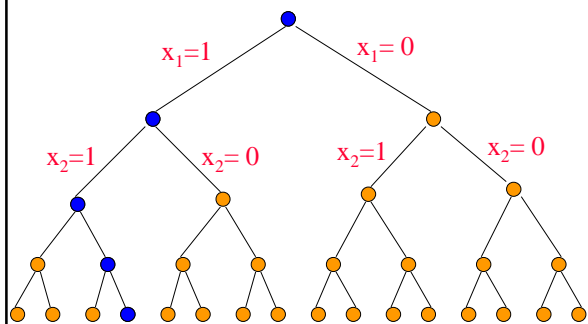
Backtracking Depth-First Search



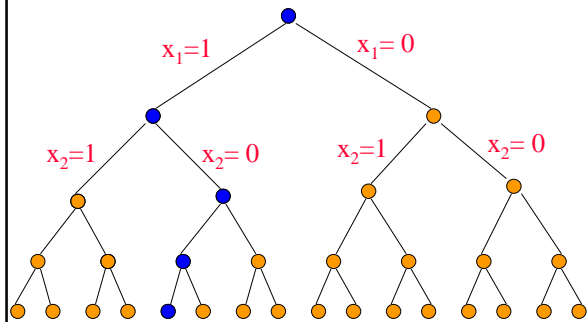
Backtracking Depth-First Search



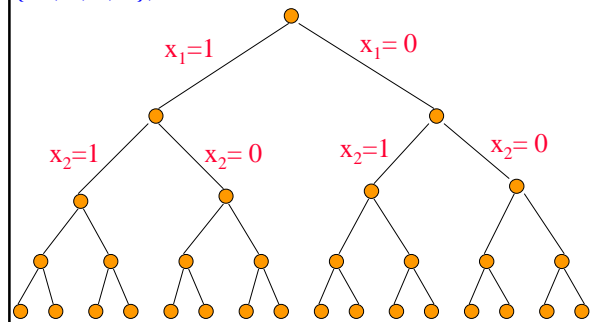
Backtracking Depth-First Search



### Backtracking Depth-First Search



### $O(2^n)$ Subset Sum & Bounding Functions {10, 5, 2, 1}, $c = 14$



Each forward and backward move takes  $O(1)$  time.

### Bounding Functions

- When a node that represents a subset whose sum equals the desired sum  $c$ , terminate.
- When a node that represents a subset whose sum exceeds the desired sum  $c$ , backtrack. I.e., do not enter its subtrees, go back to parent node.
- Keep a variable  $r$  that gives you the sum of the numbers not yet considered. When you move to a right child, check if  $\text{current subset sum} + r \geq c$ . If not, backtrack.

### Backtracking

- Space required is  $O(\text{tree height})$ .
- With effective bounding functions, large instances can often be solved.
- For some problems (e.g., 0/1 knapsack), the answer (or a very good solution) may be found quickly but a lot of additional time is needed to complete the search of the tree.
- Run backtracking for as much time as is feasible and use best solution found up to that time.

## Branch And Bound

- Search the tree using a breadth-first search (**FIFO branch and bound**).
- Search the tree as in a bfs, but replace the FIFO queue with a stack (**LIFO branch and bound**).
- Replace the FIFO queue with a priority queue (**least-cost (or max priority) branch and bound**).  
The priority of a node **p** in the queue is based on an estimate of the likelihood that the answer node is in the subtree whose root is **p**.

## Branch And Bound

- Space required is **O(number of leaves)**.
- For some problems, solutions are at different levels of the tree (e.g., 16 puzzle).

4		14	1
13	2	3	12
6	11	5	10
9	8	7	15

1	2	3	4
5	6	7	8
9	10	11	12
13	14	15	

## Branch And Bound

- FIFO branch and bound finds solution closest to root.
- Backtracking may never find a solution because tree depth is infinite (unless repeating configurations are eliminated).
- Least-cost branch and bound directs the search to parts of the space most likely to contain the answer. So it could perform better than backtracking.