

## Arrays



## 1D Array Representation In Java, C, and C++

Memory

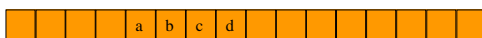


start

- 1-dimensional array  $x = [a, b, c, d]$
- map into contiguous memory locations
- $\text{location}(x[i]) = \text{start} + i$

## Space Overhead

Memory



start

space overhead = 4 bytes for **start**  
+ 4 bytes for **x.length**  
= 8 bytes  
(excludes space needed for the elements of **x**)

## 2D Arrays

The elements of a 2-dimensional array **a** declared as:

```
int [][]a = new int[3][4];
```

may be shown as a table

<b>a[0][0]</b>	<b>a[0][1]</b>	<b>a[0][2]</b>	<b>a[0][3]</b>
<b>a[1][0]</b>	<b>a[1][1]</b>	<b>a[1][2]</b>	<b>a[1][3]</b>
<b>a[2][0]</b>	<b>a[2][1]</b>	<b>a[2][2]</b>	<b>a[2][3]</b>

### Rows Of A 2D Array

a[0][0]	a[0][1]	a[0][2]	a[0][3]	→ row 0
a[1][0]	a[1][1]	a[1][2]	a[1][3]	→ row 1
a[2][0]	a[2][1]	a[2][2]	a[2][3]	→ row 2

### Columns Of A 2D Array

a[0][0]	a[0][1]	a[0][2]	a[0][3]
a[1][0]	a[1][1]	a[1][2]	a[1][3]
a[2][0]	a[2][1]	a[2][2]	a[2][3]

column 0   column 1   column 2   column 3

### 2D Array Representation In Java, C, and C++

2-dimensional array **x**

a, b, c, d

e, f, g, h

i, j, k, l

view 2D array as a 1D array of rows

**x** = [row0, row1, row 2]

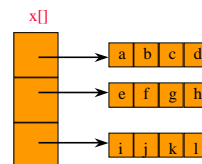
row 0 = [a,b, c, d]

row 1 = [e, f, g, h]

row 2 = [i, j, k, l]

and store as 4 1D arrays

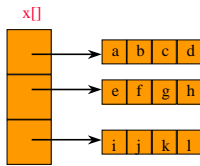
### 2D Array Representation In Java, C, and C++



**x.length** = 3

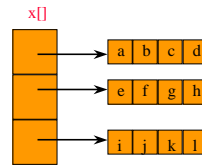
**x[0].length** = **x[1].length** = **x[2].length** = 4

## Space Overhead



space overhead = overhead for 4 1D arrays  
 $= 4 * 8 \text{ bytes}$   
 $= 32 \text{ bytes}$   
 $= (\text{number of rows} + 1) * 8 \text{ bytes}$

## Array Representation In Java, C, and C++



- This representation is called the **array-of-arrays** representation.
- Requires contiguous memory of size 3, 4, 4, and 4 for the 4 1D arrays.
- 1 memory block of size **number of rows** and **number of rows** blocks of size **number of columns**

## Row-Major Mapping

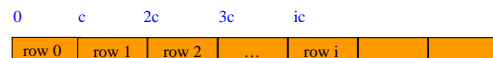
- Example 3 x 4 array:

a b c d  
e f g h  
i j k l

- Convert into 1D array  $y$  by collecting elements by rows.
- Within a row elements are collected from left to right.
- Rows are collected from top to bottom.
- We get  $y[] = \{a, b, c, d, e, f, g, h, i, j, k, l\}$



## Locating Element $x[i][j]$



- assume  $x$  has  $r$  rows and  $c$  columns
- each row has  $c$  elements
- $i$  rows to the left of row  $i$
- so  $ic$  elements to the left of  $x[i][0]$
- so  $x[i][j]$  is mapped to position  $ic + j$  of the 1D array

## Space Overhead

row 0	row 1	row 2	...	row i		
-------	-------	-------	-----	-------	--	--

4 bytes for **start** of 1D array +  
4 bytes for **length** of 1D array +  
4 bytes for **c** (number of columns)  
= 12 bytes

(number of rows = **length** / **c**)

## Disadvantage

Need contiguous memory of size **rc**.

## Column-Major Mapping

**a b c d**

**e f g h**

**i j k l**

- Convert into 1D array **y** by collecting elements by columns.
- Within a column elements are collected from top to bottom.
- Columns are collected from left to right.
- We get **y** = {**a**, **e**, **i**, **b**, **f**, **j**, **c**, **g**, **k**, **d**, **h**, **l**}

## Matrix

Table of values. Has rows and columns, but numbering begins at 1 rather than 0.

**a b c d** row 1

**e f g h** row 2

**i j k l** row 3

- Use notation **x(i,j)** rather than **x[i][j]**.
- May use a 2D array to represent a matrix.

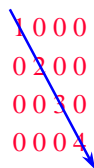
## Shortcomings Of Using A 2D Array For A Matrix

- Indexes are off by 1.
- Java arrays do not support matrix operations such as **add**, **transpose**, **multiply**, and so on.
  - Suppose that **x** and **y** are 2D arrays. Can't do **x + y**, **x - y**, **x \* y**, etc. in Java.
- Develop a class **Matrix** for object-oriented support of all matrix operations. See text.

## Diagonal Matrix

An **n x n** matrix in which all nonzero terms are on the diagonal.

## Diagonal Matrix



```

1 0 0 0
0 2 0 0
0 0 3 0
0 0 0 4
  
```

- **x(i,j)** is on diagonal iff **i = j**
- number of diagonal elements in an **n x n** matrix is **n**
- non diagonal elements are zero
- store diagonal only vs **n<sup>2</sup>** whole

## Lower Triangular Matrix

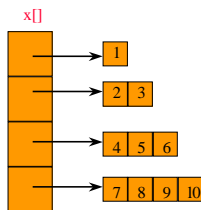
An **n x n** matrix in which all nonzero terms are either on or below the diagonal.

```

1 0 0 0
2 3 0 0
4 5 6 0
7 8 9 10
  
```

- **x(i,j)** is part of lower triangle iff **i >= j**.
- number of elements in lower triangle is **1 + 2 + ... + n = n(n+1)/2**.
- store only the lower triangle

## Array Of Arrays Representation



Use an irregular 2-D array ... length of rows is not required to be the same.

## Creating And Using An Irregular Array

```
// declare a two-dimensional array variable
// and allocate the desired number of rows
int [][] irregularArray = new int [numberOfRows][];

// now allocate space for the elements in each row
for (int i = 0; i < numberOfRows; i++)
    irregularArray[i] = new int [size[i]];

// use the array like any regular array
irregularArray[2][3] = 5;
irregularArray[4][6] = irregularArray[2][3] + 2;
irregularArray[1][1] += 3;
```

## Map Lower Triangular Array Into A 1D Array

Use row-major order, but omit terms that are not part of the lower triangle.

For the matrix

```
1 0 0 0
2 3 0 0
4 5 6 0
7 8 9 10
```

we get

1, 2, 3, 4, 5, 6, 7, 8, 9, 10

## Index Of Element [i][j]

0	1	3	6			
r1	r2	r3	...	row i		

- Order is: row 1, row 2, row 3, ...
- Row *i* is preceded by rows 1, 2, ..., *i*-1
- Size of row *i* is *i*.
- Number of elements that precede row *i* is  $1 + 2 + 3 + \dots + i-1 = i(i-1)/2$
- So element (*i*,*j*) is at position  $i(i-1)/2 + j - 1$  of the 1D array.