# **B**-splines and splines

A rough (and not fully precise) characterization of a spline is a *piecewise polynomial* that is as smooth as possible without becoming a single polynomial.

In the following, we define a basis for splines, the collection of B-basis functions, short B-splines. B-splines express a spline in B-form, i.e. as a linear combination of B-splines.

The *B*-form is particularly well-suited for applications such as graphics and geopmetric modelling, as well as numerical computations. Compared with other representations, *B*-splines have the advantage built-in continuity between the polynomial pieces – a property that is retained by linear combinations. Further desirable properties become clear as we develop the subject.

(The letter B stands for 'basic' and is denoted by capital letters, for historical reasons even though it is scalar-valued.)

### **B**-spline definition by recursion

1

Let

$$t_{(i:j)} := t_i, t_{i+1}, \dots, t_j$$
 (1)

be a nondecreasing sequence of scalars, i.e.  $t_{k+1} \ge t_k$ . The scalars are called *knots*. Then the *B*-spline of degree *d* is defined recursively as follows.

$$B(u|t_{(i:i+1)}) := \begin{cases} 1 & \text{if } t_i \leq u < t_{i+1} \\ 0 & \text{otherwise,} \end{cases}$$
$$B(u|t_{(i:i+d+1)}) := \ell(u|_{i,i+d})B(u|t_{(i:i+d)}) \\ + (1 - \ell(u|_{i+1,i+d+1}))B(u|t_{(i+1:i+d+1)}) \\ \text{where } \ell(u|_{i,j}) := \begin{cases} \frac{u-t_i}{t_j-t_i} & \text{if } t_i \neq t_j \\ 0 & \text{otherwise.} \end{cases}$$

 $X_{5-}$  Verify that for  $t_{j+d+1} > t_j$ ,  $B(u|t_{(j:j+d+1)}) > 0$  on the interval  $(t_j, t_{j+d+1})$ .

 $X_{5-}$  Show that a *B*-spline value does not change when the knots and the argument *u* are all translated by the same amount. Show that a *B*-spline value does not change when the knots and the argument are all scaled by the same amount. That is, *B*-spline are invariant under linear reparameterization.

 $X_{10}$ - We abbreviate  $\mathbf{t} := t_{(j:j+d+1)}$  and denote by  $\Pi_{d,\mathbf{t}}$  the set of all piecewise polynomial functions of degree d with breaks at the knots in  $\mathbf{t}$ . Show that  $B(u|\mathbf{t}) \in \Pi_{d,\mathbf{t}}$ .

### Splines defined

A *spline* s of degree d is a linear combination of B-splines:

$$s(u|t_{(i:i+d+n+1)}) := \sum_{j=i}^{i+n} c_{j:j+d+1} B(u|t_{(j:j+d+1)}), \quad c_{j:k} \in \mathbb{R}.$$
 (2)

#### Jorg's Lecture Notes – B-splines and splines

We note that there are n+d+1 knots for n coefficients. We sometimes abbreviate  $c_j := c_{j:j+d+1}$  and  $B_j := B(u|t_{(j:j+d+1)})$  when the degree d is understood.

We will see that we can evaluate the spline on the interval  $[t_{i+d}, t_{i+n+1}]$ . That is, we will need d additional knots on each side of the interval. A knot sequence

$$\mathbf{t}_Z := \{\ldots, -2, -1, 0, 1, 2, \ldots\}$$

gives rise to the *uniform* splines (also called *Cardinal Splines*):

$$B(u|i:i+d+1)) = B(u-i|0:d+1)).$$

The recursion for uniform splines simplifies to

$$B(u|0:d+1) = uB(u|0:d) + (d+1-u)B(u|1:d+1).$$

The knot sequence

$$\mathbf{t}_B := \{\ldots, 0, 0, 0, 1, 1, 1, \ldots\}$$

with  $\mu + 1$  zeros and  $\nu + 1$  ones yields polynomials in Bernstein-Bezier form

$$B_{\mu,\nu}(u) := \binom{\mu+\nu}{\mu} (1-u)^{\mu} u^{\nu} = u B_{\mu,\nu-1}(u) + (1-u) B_{\mu-1,\nu}(u)$$

#### **Spline Evaluation**

We can be more efficient than evaluating each *B*-spline separately and forming linear combinations. The algorithm for obtaining the value directly from the coefficients and the knots is called *de Boor algorithm*. In the special case  $t = t_B$  the de Boor algorithm is also called *de Casteljau's* algorithm.

The key to evaluating a spline by recursion is to express coefficients at level d+1 as a linear polynomial in u of two coefficients at level d:

$$c_{i:i+d} := (1 - \ell(u|_{i,i+d})) c_{i-1:i+d} + \ell(u|_{i,i+d}) c_{i:i+d+1}$$
(3)  
$$=^* \frac{t_{i+d} - u}{t_{i+d} - t_i} c_{i-1:i+d} + \frac{u - t_i}{t_{i+d} - t_i} c_{i:i+d+1}.$$

Here  $=^*$  indicates the case that  $t_{i+d} \neq t_i$ . While, formally, all  $c_{i:j}$  are functions of u, in the following context u will be the parameter of evaluation, and therefore a fixed number so that the expressions in (3) will be constants.

To evaluate the spline  $s(u|t_{(i:i+d+n+1)})$  at a value

$$u \in [t_j, t_{j+1}) \subset [t_{i+d}, t_{i+n+1}),$$

we compute  $c_{j:j+1}$  by repeatedly applying Equation (3) (cf. ). It is convenient



Figure 1: The indices of splines obtained by recursion upwards, i.e. (a:b) for indices a and b corresponds to  $B(u|t_{(a:b)})$  and  $c_{a:b}$ . For the reverse recursion for evaluation, the weights for forming (a:b) are  $1 - \ell(u|_{a:b}) = \frac{t_b - u}{t_b - t_a}$  and  $\ell(u|_{a:b}) = \frac{u - t_a}{t_b - t_a}$  respectively.

to look at a window of relevant knots  $t_{j-d:j+d+1}$ .

$$f(u|t_{(j-d:j+d+1)}) = \sum_{i=j}^{j+d} c_{i-d:i+1} B(u|t_{(i-d:i+1)})$$
  
= 
$$\sum_{i=j}^{j+d-1} c_{i-d+1:i+1} B(u|t_{(i-d+1:i+1)}) = \dots = \sum_{i=j}^{j} c_{i:i+1} B(u|t_{(i:i+1)})$$
  
= 
$$c_{j:j+1}$$

The last equality follows since splines with two knots are piecewise constant and for  $u \in [t_j, t_{j+1})$ ,  $B(u|t_{(i:i+1)}) = 0$  except that it is 1 when i = j.

Example: We evaluate at  $u = 0 \in [t_j, t_{j+1}]$ , a spline defined by d = 2 and

$$\mathbf{t} = \begin{bmatrix} -3 & -2 & -1 & 1 & 5 & 6 \end{bmatrix}, \quad c = \begin{bmatrix} 48 & 12 & 24 \end{bmatrix}$$

Therefore  $t_j = -1$  and  $c_{j-2:j+1} = 48$ . We compute

$$\ell(u|_{j-1:j+1}) = \frac{u - t_{j-1}}{t_{j+1} - t_{j-1}} = \frac{0 - (-2)}{1 - (-2)} = \frac{2}{3}$$
$$\ell(u|_{j:j+2}) = \frac{u - t_j}{t_{j+2} - t_j} = \frac{0 - (-1)}{5 - (-1)} = \frac{1}{6}$$
$$\ell(u|_{j:j+1}) = \frac{u - t_j}{t_{j+1} - t_j} = \frac{0 - (-1)}{1 - (-1)} = \frac{1}{2}$$

The schema then reads

$c_{j-2:j+1}:$	48		12		24
$\ell(u _{j:j+2}):$		2/3		1/6	
$c_{j-1:j+1}:$		24		14	
$\ell(u _{j:j+1}):$			1/2		
$c_{j:j+1}$ :			19		

In the example, we need 6 knots and 3 coefficients to have one interval, [-1..1]. The define the three degree 2 B-splines that have support (are non-zero) on the interval.

 $X_5$ – Using the data of the example, associate each coefficient  $c_j$  with its *Greville* abscissa defined by

$$t_j^* := \frac{1}{d} \sum_{i=1}^d t_{j+i}$$

and draw the control polygon and the de Boor evaluation.  $X_{3-}$  Show that for  $t_{j} \leq u \leq t_{j+1}$  all  $0 \leq \ell(u|_{i:j}) \leq 1$ .  $X_{3-}$  Check that the algorithm is well-defined for multiple knots.

# Differentiation

The spline  $s(u|t_{(i:i+d+n+1)}) := \sum_{j=i}^{i+n} c_{j:j+d+1} B(u|t_{(j:j+d+1)})$  of degree d has the derivative

$$s'(u|t_{(i:i+d+n+1)}) := \sum_{j=i+1}^{i+n} c'_{j:j+d} B(u|t_{(j:j+d)}), \tag{4}$$

where

$$c'_{j:j+d} := d \ \frac{c_{j:j+d+1} - c_{j-1:j+d}}{t_{j+d} - t_j}$$

#### Multiplication

 $X_{20}$ - What challenges do you see in deriving the coefficients and knots of the spline that is the product of two given splines? (Conversion to BB-form and back helps, but here the structure is asked)

## Integration

The spline  $s(u|t_{(i:i+d+n-1)}) := \sum_{j=i}^{i+n} c_{j:j+d} B(u|t_{(j:j+d)})$ , has the antiderivative

$$\sum_{j=i}^{i+n} c_{j:j+d+1}^* B(u|t_{(j:j+d+1)}), \qquad c_{j:j+d+1}^* := \text{const} + \begin{cases} \sum_{i=j_0}^j c_{i:i+d} \frac{t_{i+d}-t_i}{d} \\ \sum_{i=j}^{j_0-1} c_{i:i+d} \frac{t_{i+d}-t_i}{d} \end{cases}$$
(5)

where  $j_0$  and the constant are arbitrary ( $j_0$  separates two subsets of indices).

## Continuity

The knots  $t_j$  and  $t_{j+1}$  may agree. Such coalescing of knots decreases the otherwise guaranteed continuity of the B-spline and hence of the spline  $s_{d,t}$  as follows. If m is the multiplicity of the knot  $t_j$ , i.e.  $\ldots < t_j = \ldots = t_{j+m-1} < \ldots$  then  $s_{d,\mathbf{t}}$  is at least

$$k = d - m$$

times continuously differentiable at  $t_i$ .

## Reproduction

Marsden's identity

$$(u-\tau)^{d} = \sum_{j=0}^{d} a_{j:j+d}(\tau) B(u|\mathbf{t}), \quad \mathbf{t} := t_{(j:j+d+1)},$$
$$a_{j:j+d}(\tau) := (t_{j+1}-\tau) \cdots (t_{j+d}-\tau)$$

shows what coefficients  $c_t$  the spline  $\sum_j c_t B(u|t)$  has to have so that so that the spline reproduces a given polynomial  $p := \sum_{j=0}^{k} (u-\tau)^{j}$ . Differentiating *i* times with respect to  $\tau$  and dividing by *d*! yields

$$\frac{(u-\tau)^{d-j}}{(d-j)!} = \frac{(-1)^j}{d!} \sum_i B(u|t_{(i:I+d+1)}) D^j a_{i:i+d}(\tau).$$

Taylor expansion gives the coefficients,

$$c_i = \lambda(d, i)p := \sum_{j=0}^d \frac{(-1)^j (D^j a_{i:i+d})(\tau)}{d!} (D^{d-j}p)(\tau).$$

The  $\lambda(d, i)$  are called *dual functionals*.

 $X_{5}$ - Show that for fixed degree d and any knot sequence  $(t_i)_{i=-\infty..\infty}$ 

$$\sum_{j} B(u|t_{(j:j+d+1)}) = 1.$$
(6)

 $X_{5}$ - Show for any linear function p and Greville abscissae  $t_{i}^{*}$ 

$$\sum_{j} B(u|t_{(j:j+d+1)})p(t_{j}^{*}) = p.$$
(7)

 $X_{10}$ - Show that for  $\tau \in [x_j, x_{j+d+1})$ 

$$\lambda(d,i)(\sum_{j} B(u|t_{(j:j+d+1)})c_j) = c_i.$$
(8)

Of course splines are not usually used to reproduce polynomials but to approximate (or even interpolate) a sequence of points.

# Interpolation

(c.f. cubic spline interpolation)

The Schoenberg-Whittney interpolation theorem: given a vector of interpolation points  $u_j$  and values  $f(u_j)$ , there always exists a spline of degree d that is d-1 times continuously differentiable and interpolates the data. A necessary and sufficient condition for the banded system of interpolation constraints to be solvable is  $t_j < u_j < t_{j+d+1}$ .

### Stability and local well-conditioning of a spline basis

Small spline function implies small spline coefficients  $(c_j \neq 0 \text{ for some } j)$ :

$$2^{-d} \max_{i} \|c_i\| < \|\sum c_j B_j\|.$$

 $\|\sum c_j B_j\|_{\infty} \le \max_i \|c_i\|$  follows from  $\sum B_j = 1$ .

# Knot insertion and subdivision.

Let  $C_{c,t}$  be the control polygon of the spline s. Then

$$\sup_{u} |s(u|t) - C_{c,t}(u)| \le const \sup_{u} |D^2 s(u)| |\sup_{i} (t_{j+1} - t_j)|^2.$$

That is, the control polygon *converges with quadratic error* to s as the distance between the knots decreases. Inserting a knot requires generating new coefficients without changing the spline. The correct way to do this when the knot  $\hat{t}$  is inserted into the sequence t is to choose the *new coefficient sequence* as follows.

$$\hat{c}_j := \begin{cases} c_j, & t_{j+d} \le \hat{t}; \\ (1 - \ell_{d,j}(\hat{t}))c_{j-1} + \ell_{d,j}(\hat{t})c_j, & t_j < \hat{t} < t_{j+d}; \\ c_{j-1}, & \hat{t} \le t_j \end{cases}$$

Example: Given the spline defined by

$$t = \begin{bmatrix} -3 & -2 & -1 & 1 & 5 & 6 \end{bmatrix},$$
  
$$c = \begin{bmatrix} 48 & 12 & 24 \end{bmatrix}.$$

insert a new knot  $\hat{t} = 0$ . We compute

$$\begin{array}{cccccccc} c & 48 & 12 & 24 \\ w_2: & 2/3 & 1/6 \\ \hat{c}: & 48 & 24 & 14 & 24 \end{array}$$

The *Greville abscissae* are

 $t^* = -1.5, 0, 3,$  before insertion and  $t^* = -1.5, -.5, .5, 3,$  after insertion.

 $X_{10}$ - Given the spline *curve* defined by

$$t = \begin{bmatrix} 0, 0, 0, 0, 1, 3, 5, 5, 5, 5 \end{bmatrix},$$
  
$$c = \begin{bmatrix} -1 & 0 & 1 & 2 & 3 & 4 \\ 0 & 0 & 1 & 1 & 0 & 1 \end{bmatrix}$$

compute the new control polygon after insertion of the new knot  $\hat{t} := 2$ . (Hints: there are knots  $x_0, ..., x_9$  and 6 intervals, hence 9 = 6 + d. For a curve we do not display the Greville abscissae but only the (x, y) pairs.)

 $X_{10}$ - Given the spline function of degree 3 with coefficients c and  $t = t_{2Z} = \dots, -2, 0, 2, \dots$  give the coefficient sequence after insertion of knots at every odd integer; that is after changing the knot sequence to  $t_Z$ .

### The convex hull property

 $X_{3-}$  Show that for  $t_j \leq u < t_{j+1}$ , s(u) is a convex combination of  $c_{j-d}, \ldots, c_j$ .

#### Variation diminuition

The number of strong sign changes in any increasing sequence of values  $p_d(x_1), \ldots, p_d(x_r)$  of the spline is less than the number of strong sign changes in the sequence of coefficients  $c_i$ . This is often written as

$$S^{-}(s) \le S^{-}(c),$$

or  $S^{-}(s(t_1), \ldots, s(t_n)) \leq S^{-}(c_1, \ldots, c_n)$  for t strictly increasing. Proof: insert  $t_i$  into the knot sequence to obtain new coefficients by convex averaging.

#### Shape preservation

A spline crosses any straight line no more often than its control polygon. Proof:  $S(s-\ell) \leq S(c-\ell)$ .

 $X_{3-}$  Show that shape preservation implies that a spline is *monotone* if its control polygon is monotone.  $X_{3-}$  Show that shape preservation implies that a spline is *convex* if its control polygon is convex.

## Alternative definitions of B-splines

(that allow the definition of splines in several dimensions)

1. Hermite-Genocchi formula: B-spline as distribution.

$$\int_{R} B_{d}(x, u)\phi(u)du$$
$$= d! \int_{\Delta(d)} \phi(\lambda x)d\lambda$$

where  $\Delta(d)$  is the *d*-simplex spanned by  $e_0, e_1, \ldots, e_d, e_0 = 0, e_i$  the *i*th unit vector.

2. Geometric interpretation. With T a d-simplex such that the first component of each vertex is one of the knots  $x_j$  and  $\pi$  the projection  $\mathbb{R}^d \mapsto \mathbb{R}$ :  $\pi(x_1, x_2, \ldots, x_d) = x_1$ . Then

$$B_d(x, u) := \operatorname{vol}_{d-1}(T \cap (\pi^{-1}(u))) / \operatorname{vol}_d(T).$$