

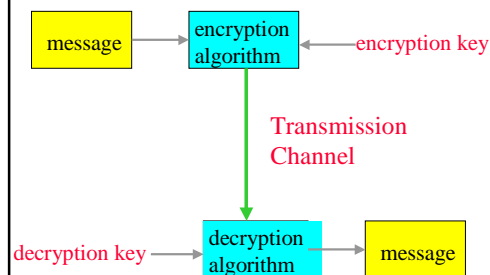


## Hard Problems



- Some problems are hard to solve.
  - No polynomial time algorithm is known.
  - E.g., NP-hard problems such as machine scheduling, bin packing, 0/1 knapsack.
- Is this necessarily bad?
- Data encryption relies on difficult to solve problems.

## Cryptography



## Public Key Cryptosystem (RSA) 🗝️

- A public encryption method that relies on a public encryption algorithm, a public decryption algorithm, and a public encryption key.
- Using the public key and encryption algorithm, everyone can encrypt a message.
- The decryption key is known only to authorized parties.
- Asymmetric method.
  - Encryption and decryption keys are different; one is not easily computed from the other.

## Public Key Cryptosystem (RSA) 🗝️

- $p$  and  $q$  are two prime numbers.
- $n = pq$
- $m = (p-1)(q-1)$
- $a$  is such that  $1 < a < m$  and  $\gcd(m,a) = 1$ .
- $b$  is such that  $(ab) \bmod m = 1$ .
- $a$  is computed by generating random positive integers and testing  $\gcd(m,a) = 1$  using the extended Euclid's gcd algorithm.
- The extended Euclid's gcd algorithm also computes  $b$  when  $\gcd(m,a) = 1$ .

## RSA Encryption And Decryption 🗝️

- Message  $M < n$ .
- Encryption key =  $(a,n)$ .
- Decryption key =  $(b,n)$ .
- Encrypt  $\Rightarrow E = M^a \bmod n$ .
- Decrypt  $\Rightarrow M = E^b \bmod n$ .

## Breaking RSA 🗝️

- Factor  $n$  and determine  $p$  and  $q$ ,  $n = pq$ .
- Now determine  $m = (p-1)(q-1)$ .
- Now use Euclid's extended gcd algorithm to compute  $\gcd(m,a)$ .  $b$  is obtained as a byproduct.
- The decryption key  $(b,n)$  has been determined!

## Security Of RSA



- Relies on the fact that prime factorization is computationally very hard.
- Let  $q$  be the number of bits in the binary representation of  $n$ .
- No algorithm, polynomial in  $q$ , is known to find the prime factors of  $n$ .
- Try to find the factors of a 100 bit number.

## Elliptic Curve Cryptography (ECC)

- Asymmetric Encryption Method
  - Encryption and decryption keys are different; one is not easily computed from the other.
- Relies on difficulty of computing the discrete logarithm problem for the group of an elliptic curve over some finite field.
  - Galois field of size a power of 2.
  - Integers modulo a prime.
- 1024-bit RSA ~ 200-bit ECC (cracking difficulty).
- Faster to compute than RSA?

## Data Encryption Standard

- Used for password encryption.
- Encryption and decryption keys are the same, and are secret.
- Relies on the computational difficulty of the satisfiability problem.
- The satisfiability problem is NP-hard.

## Satisfiability Problem

- The permissible values of a boolean variable are **true** and **false**.
- The **complement** of a boolean variable  $x$  is denoted  $\bar{x}$ .
- A **literal** is a boolean variable or the complement of a boolean variable.
- A **clause** is the **logical or** of two or more literals.
- Let  $x_1, x_2, x_3, \dots, x_n$  be  $n$  boolean variables.

## Satisfiability Problem

- Example clauses:
  - $x_1 + x_2 + x_3$
  - $x_4 + \bar{x}_7 + x_8$
  - $x_3 + x_7 + x_9 + x_{15}$
  - $x_2 + x_5$
- A **boolean formula** (in conjunctive normal form CNF) is the **logical and** of  $m$  clauses.
- $F = C_1 C_2 C_3 \dots C_m$

## Satisfiability Problem

- $F = (x_1 + x_2 + x_3)(\bar{x}_4 + \bar{x}_7 + x_8)(x_2 + x_5)$
- $F$  is **true** when  $x_1, x_2$ , and  $x_4$  (for e.g.) are **true**.

## Satisfiability Problem

- A boolean formula is **satisfiable** iff there is at least one truth assignment to its variables for which the formula evaluates to **true**.
- Determining whether a boolean formula in CNF is satisfiable is NP-hard.
- Problem is solvable in polynomial time when no clause has more than **2** literals.
- Remains NP-hard even when no clause has more than **3** literals.

## Other Problems

- Partition
  - Partition **n** positive integers  $s_1, s_2, s_3, \dots, s_n$  into two groups **A** and **B** such that the sum of the numbers in each group is the same.
  - **[9, 4, 6, 3, 5, 1, 8]**
  - **A = [9, 4, 5]** and **B = [6, 3, 1, 8]**
- NP-hard.

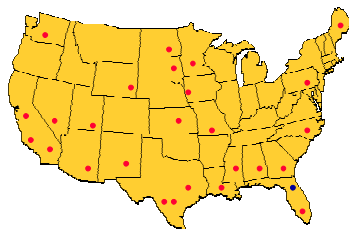
## Subset Sum Problem

- Does any subset of **n** positive integers  $s_1, s_2, s_3, \dots, s_n$  have a sum exactly equal to **c**?
- **[9, 4, 6, 3, 5, 1, 8]** and **c = 18**
- **A = [9, 4, 5]**
- NP-hard.

## Traveling Salesperson Problem (TSP)

- Let **G** be a weighted directed graph.
- A **tour** in **G** is a cycle that includes every vertex of the graph.
- **TSP**  $\Rightarrow$  Find a tour of shortest length.
- Problem is NP-hard.

## Applications Of TSP



- Home city
- Visit city

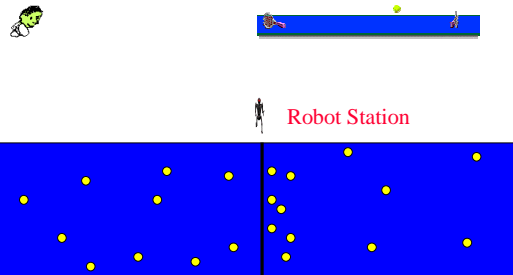
## Applications Of TSP

- Each vertex represents a city that is in Joe's sales district.
- The weight on edge **(u,v)** is the time it takes Joe to travel from city **u** to city **v**.
- Once a month, Joe leaves his home city, visits all cities in his district, and returns home.
- The total time he spends on this tour of his district is the travel time plus the time spent at the cities.
- To minimize total time, Joe must use a shortest-length tour.

## Applications Of TSP

- Tennis practice.
- Start with a basket of approximately 200 tennis balls.
- When balls are depleted, we have 200 balls lying on and around the court.
- The balls are to be picked up by a robot (more realistically, the tennis player).
- The robot starts from its station visits each ball exactly once (i.e., picks up each ball) and returns to its station.

## Applications Of TSP

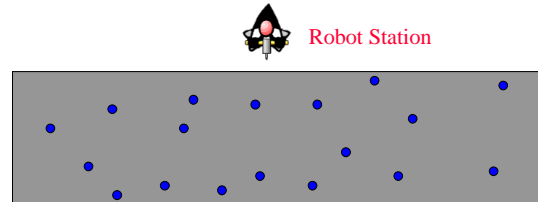


## Applications Of TSP

- 201 vertex TSP.
- 200 tennis balls and robot station are the vertices.
- Complete directed graph.
- Length of an edge  $(u,v)$  is the distance between the two objects represented by vertices  $u$  and  $v$ .
- Shortest-length tour minimizes ball pick up time.
- Actually, we may want to minimize the sum of the time needed to compute a tour and the time spent picking up balls using the computed tour.

## Applications Of TSP

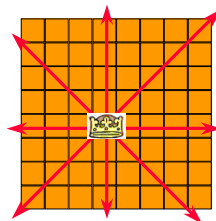
- Manufacturing.
- A robot arm is used to drill  $n$  holes in a metal sheet.



$n+1$  vertex TSP.

## n-Queens Problem

A queen that is placed on an  $n \times n$  chessboard, may attack any piece placed in the same column, row, or diagonal.



8x8 Chessboard

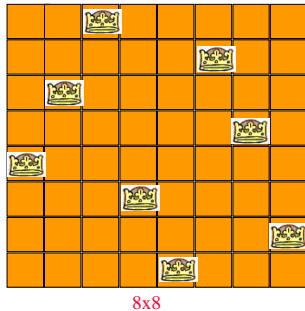
## n-Queens Problem

Can  $n$  queens be placed on an  $n \times n$  chessboard so that no queen may attack another queen?



4x4

## ♔ n-Queens Problem ♚



8x8

## Difficult Problems

- Many require you to find either a **subset** or **permutation** that satisfies some constraints and (possibly also) optimizes some objective function.
- May be solved by organizing the **solution space** into a **tree** and **systematically searching** this tree for the answer.

## Subset Problems

- Solution requires you to find a **subset** of **n** elements.
- The subset must satisfy some constraints and possibly optimize some objective function.
- Examples.
  - Partition.
  - Subset sum.
  - 0/1 Knapsack.
  - Satisfiability (find subset of variables to be set to **true** so that formula evaluates to **true**).
  - Scheduling **2** machines.
  - Packing **2** bins.

## Permutation Problems

- Solution requires you to find a **permutation** of **n** elements.
- The **permutation** must satisfy some constraints and possibly optimize some objective function.
- Examples.
  - TSP.
  - n-queens.
    - Each queen must be placed in a different row and different column.
    - Let queen **i** be the queen that is going to be placed in row **i**.
    - Let  $c_i$  be the column in which queen **i** is placed.
    - $c_1, c_2, c_3, \dots, c_n$  is a permutation of  $[1, 2, 3, \dots, n]$  such that no two queens attack.

## Solution Space

- Set that includes at least one solution to the problem.
- Subset problem.
  - $n = 2, \{00, 01, 10, 11\}$
  - $n = 3, \{000, 001, 010, 100, 011, 101, 110, 111\}$
- Solution space for subset problem has  $2^n$  members.
- Nonsystematic search of the space for the answer takes  $O(p2^n)$  time, where **p** is the time needed to evaluate each member of the solution space.

## Solution Space

- Permutation problem.
  - $n = 2, \{12, 21\}$
  - $n = 3, \{123, 132, 213, 231, 312, 321\}$
- Solution space for a permutation problem has **n!** members.
- Nonsystematic search of the space for the answer takes  $O(pn!)$  time, where **p** is the time needed to evaluate a member of the solution space.